### QUASI-ORDINARY SINGULARITIES AND NEWTON TREES

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ABSTRACT. In this paper we study some properties of the class of  $\nu$ -quasi-ordinary hypersurface singularities. They are defined by a very mild conditions on their (projected) Newton polygon. We associate with them a Newton tree and characterize quasi-ordinary hypersurface singularities among  $\nu$ -quasi-ordinary hypersurface singularities in terms of their Newton tree. A formula to compute the discriminant of a quasi-ordinary Weierstrass polynomial in terms of the decorations of its Newton tree is given. This allows to compute the discriminant avoiding the use of determinants and even for not Weierstrass prepared polynomials. This is important for applications like algorithmic resolutions. We compare the Newton tree of a quasi-ordinary singularity and those of its curve transversal sections. We show that the Newton trees of the transversal sections do not give the tree of the quasi-ordinary singularity in general. They do if we know that the Newton tree of the quasi-ordinary singularity has only one arrow.

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#### Introduction

Let  $\mathbb{K}$  be an algebraically closed field of characteristic zero. Let  $f \in \mathbb{K}[[x]][z]$  be a polynomial with coefficients in the formal power series ring  $\mathbb{K}[[x]] := \mathbb{K}[[x_1, \ldots, x_d]]$ . If d = 1, its zero locus ( $\{f = 0\}, 0$ )  $\subset (\mathbb{K}^2, 0)$  defines a germ of plane curve singularity. One of the initial ideas to deal with germs of plane curve singularities is the use of the Newton algorithm. Inside the algorithm, the first step is governed by the Newton polygon of f and, for each compact face of the Newton polygon, one proceeds by doing the Newton process, and so on. This procedure is codified in a tree called Newton tree. Many properties can be codified in the tree, for instance, using bi-colored Newton trees one can compute the intersection multiplicity of two plane curves. This allows to compute from the Newton tree the Milnor number

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of a germ since it can be expressed in terms of intersection multiplicities. More recently for a finite codimension ideal I in  $\mathbb{K}[[x_1, x_2]]$  its multiplicity, its Lojasiewicz exponent, and its Hilbert–Samuel multiplicity can be computed from its Newton tree, see [10].

For d > 1 the well-known class of hypersurface singularities which generalizes the case of curves is the class of quasi-ordinary hypersurface singularities. In that class, one of the key results is a factorization theorem given by Jung=-Abhyankar, see [1]. One feature of quasi-ordinary singularities is that in some coordinates their Newton polyhedron is a polygonal path, i.e., all its compact faces have dimension one. Then we can apply the Newton process as in the case of curves. Moreover, after Newton procedure, the condition to be quasi-ordinary is preserved. In [5], we have studied a generalization of quasi-ordinary singularities, called  $\nu$ -quasi-ordinary by H. Hironaka [19, Definition 6.1]. Roughly speaking for a  $\nu$ -quasi-ordinary hypersurface defined by a polynomial  $f \in \mathbb{K}[[x]][z]$  we only require the upper part of the Newton polyhedron of f to be a polygonal path. We apply the Newton process associated with these faces of dimension 1 and we iterate the process whenever the upper part of the corresponding Newton polyhedron is a polygonal path. As in the case of curves we encode the Newton process in a tree, but the tree bears leaves (arrows) and fruits (black boxes). Using this tree, we describe the condition for two  $\nu$ -quasi-ordinary series f and  $g \in \mathbb{K}[[x]][z]$  to have z-resultant which is a monomial times a unit. In particular, a quasi-ordinary singularity has a Newton tree with only arrows.

The first main result of this article is that a  $\nu$ -quasi-ordinary polynomial f is a quasi-ordinary if and only if there exists a suitable system of coordinates for f such that its Newton tree has only arrow-heads decorated with (0) and (1) and has no black boxes, see Theorem 4.4. This result can also be stated as follows:

Let f be a  $\nu$ -quasi-ordinary polynomial. Then, f is a quasi-ordinary polynomial if and only if f is reduced,  $\nu$ -quasi-ordinary and the successive transforms by Newton maps are also  $\nu$ -quasi-ordinary.

This gives a characterization of quasi-ordinary polynomial which does not depend on the definition of Newton trees, that is, quasi-ordinary is equivalent to reduced and stably (by Newton maps)  $\nu$ -quasi-ordinary. Theorem 4.4 is proved studying how the derivative  $\partial f/\partial z$  separates from f on the Newton tree of the product. In the case of curves this is an algebraic elementary way to recover Lê-Michel-Weber theorem [22]. We also give the formula to compute the discriminant of a quasi-ordinary Weierstrass polynomial f in terms of the decorations of the Newton tree of f, see Proposition 4.17. This formula has two important consequences. The first is that this formula allows to compute the discriminant avoiding the use of determinants. The second one, if one starts with a not Weierstrass polynomial, the computation of the Weierstrass decomposition is not needed. This is important for applications, for instance algorithmic resolutions, see [6], [7], [8]. In [6] the author gives an effective algorithm for desingularization of surfaces using Jung's method. His Algorithms 2 and 3 use as input a monic polynomial. If the polynomial is not monic, it is a difficult task to make effective the Weierstrass preparation theorem.

Using our method and proceeding as in [2] one can avoid this problem and adapt Algorithms 2 and 3 without the monic condition.

The second main result is contained in Section 5, where we compare the Newton tree of a quasi-ordinary singularity and those of its curve transversal sections. Examples 5.8 and 5.10 show that, in general, it is not possible to recover the Newton tree of the quasi-ordinary singularity from the Newton trees of the curve transversal sections. It does if we know that the Newton tree of the quasi-ordinary singularity has only one arrow, see Theorem 5.12. However the global decorations which appear on the Newton tree of the quasi-ordinary singularities are those which appear on the Newton tree of the transversal sections. This fact plays a crucial role in the proof of the monodromy conjecture for quasi-ordinary singularities [4]. Moreover in [15] a description of the motivic Milnor fibre of an irreducible quasi-ordinary polynomial is given proving that it is a topological invariant. In [17] the notion of linear Newton tree for an irreducible quasi-ordinary polynomial is introduced and some of its properties are discussed, as its normalization or its semigroup.

In Section 1 we recall the definitions and some properties of  $\nu$ -quasi-ordinary polynomials and Newton maps. The construction of the Newton trees introduced in [5] is recalled and explained in Section 2; the notion of comparable polynomials and the computation of resultants are also recalled in this section. In order to discuss the unicity of Newton trees we work with the notion of P-good coordinates in Section 3, introduced by P. González Pérez in [14]. The last two sections are devoted to state and prove the main results of the paper.

### 1. On $\nu$ -Quasi-Ordinary Polynomials and Newton Maps

1.1. Basic facts on Newton polyhedra. We shall follow the terminology of [4], [5]. The d-tuples will be denoted in bold letters, e.g.,  $x := (x_1, \ldots, x_d)$ , and we will use the following notations:

- $\boldsymbol{x}^{\boldsymbol{\alpha}} := x_1^{\alpha_1} \cdot \ldots \cdot x_d^{\alpha_d};$   $\boldsymbol{p} \cdot \boldsymbol{q} := p_1 q_1 + \cdots + p_d q_d;$   $\boldsymbol{p} * \boldsymbol{q} := (p_1 q_1, \ldots, p_d q_d);$

- $\bullet \ \frac{\boldsymbol{p}}{\boldsymbol{q}} := \left(\frac{p_1}{q_1}, \dots, \frac{p_d}{q_d}\right).$

For Newton theory, we shall also follow the terminology of [3], [20], [23]; note that for the terms polyhedron and diagram we follow the convention in [4], [5], [23]; the terms are exchanged in [3]. Let  $\mathbb{N} \subset \mathbb{R}_+$  be the sets of non-negative integers and non-negative real numbers respectively. Let  $E \subset \mathbb{N}^{d+1}$  be a set of points and  $d \geqslant 1$ .

- The Newton diagram  $\mathcal{N}_+(E)$  is defined by the convex hull in  $(\mathbb{R}_+)^{d+1}$  of the set  $E + (\mathbb{R}_+)^{d+1}$ .
- The Newton polyhedron  $\mathcal{N}(E)$  of E is defined by the union of all compact faces of the Newton diagram of E.
- The smallest set  $E_0$  such that  $\mathcal{N}_+(E_0) = \mathcal{N}_+(E)$  is called the set of vertices

 A diagram is called polygonal if the maximal dimension of its compact faces is one.

Let  $\mathbb{K}$  be an algebraically closed field of characteristic zero. We write  $f(y) = f(y_1, \ldots, y_{d+1})$  for a formal power series of several variables  $\mathbb{K}[[y]]$ . If the Taylor expansion of f(y) is  $\sum_{\alpha \in \mathbb{N}^{d+1}} c_{\alpha} y^{\alpha}$ , then the *support* Supp(f) is defined to be  $\{\alpha \in \mathbb{N}^{d+1} : c_{\alpha} \neq 0\}$ .

- The Newton polyhedron  $\mathcal{N}(f)$  of f (resp. the Newton diagram  $\mathcal{N}_+(f)$ ) of f is defined by the Newton polyhedron (resp. the Newton diagram) of the set  $\operatorname{Supp}(f)$ .
- If  $\gamma$  is a compact face of  $\mathcal{N}_+(f)$  then the (weighted-homogeneous polynomial)  $f_{\gamma}(\boldsymbol{y}) := \sum_{(\boldsymbol{\alpha}) \in \gamma} c_{\boldsymbol{\alpha}} \boldsymbol{y}^{\boldsymbol{\alpha}} \in \mathbb{K}[\boldsymbol{y}]$  is called the *polynomial associated* with  $\gamma$ .
- If the Newton diagram  $\mathcal{N}_+(f)$  is polygonal we also say that the polyhedron  $\mathcal{N}(f)$  is a monotone polygonal path.

If  $f = f_1 \cdot \ldots \cdot f_s$  in  $\mathbb{K}[[\boldsymbol{y}]]$  then

$$\mathcal{N}_+(f) = \mathcal{N}_+(f_1) + \ldots + \mathcal{N}_+(f_s),$$

where the sum is the Minkowski sum of Newton diagrams. This implies the following result which is well-known for experts, see for example [12, Lemma 12].

**Lemma 1.1.** If  $f \in \mathbb{K}[[y]]$  has a Newton polyhedron which is a monotone polygonal path, then any irreducible factor of f which is not associated with  $y_i$ ,  $i = 1 \dots, d+1$ , has a Newton polyhedron which is a monotone polygonal path.

Let  $\gamma$  be a 1-dimensional compact face of  $\mathcal{N}(f)$ . There exist two integral points A,  $A_1$  in  $\mathbb{N}^{d+1}$  such that  $\gamma$  is the edge  $[A, A_1]$ . Let c denote the greatest common divisor of all coordinates of the vector  $A_1 - A$  and let  $\mathbf{u} := \frac{1}{c}(A_1 - A) \in \mathbb{Z}^{d+1}$ . The number of points with integer coordinates on  $\gamma$  is c+1 and any of them is of the form  $A_j = A + j\mathbf{u}$ , with  $0 \le j \le c$ . Then

$$f_{\gamma}(\boldsymbol{y}) = \sum_{a \in \gamma} c_a \boldsymbol{y}^a = \boldsymbol{y}^A \left( \sum_{j=0}^c c_{A_j} \boldsymbol{y}^{ju} \right) = \boldsymbol{y}^A p(f, \gamma)(\boldsymbol{y}^u), \tag{1.1}$$

where  $p(f, \gamma)(t)$  is the polynomial  $p(f, \gamma)(t) := \sum_{j=0}^{c} c_{A_j} t^j$  of degree c; since A and  $A_1$  are vertices of  $\gamma$  we have  $p(f, \gamma)(0) \neq 0$ . Let us factor  $p(f, \gamma)(t)$  as

$$p(f, \gamma)(t) = a_{\gamma} \prod_{i=1}^{\infty} (t - \mu_{i})^{m_{i}},$$

$$\sum_{i=1}^{\infty} m_{i} = c, \quad \mu_{i} \in \mathbb{K}^{*}, \quad \mu_{i} \neq \mu_{j}, \quad \text{if } i \neq j, \quad \text{and } a_{\gamma} \in \mathbb{K}^{*}.$$

$$(1.2)$$

Let us summarize these facts.

**Proposition 1.2.** Let  $\gamma = [A, A_1]$  be a 1-dimensional compact face of  $\mathcal{N}(f)$ . With the notations of (1.1) and (1.2) the polynomial  $f_{\gamma}(\mathbf{y})$  decomposes in  $\mathbb{K}((\mathbf{y}))$  as  $a_{\gamma}\mathbf{y}^{A}\prod(\mathbf{y}^{u}-\mu_{i})^{m_{i}}$ .

**Theorem 1.3** [12, Theorem 3]. If f in  $\mathbb{K}[[y]]$  is irreducible and has a Newton polyhedron  $\mathcal{N}(f)$  which is a monotone polygonal path, then the diagram  $\mathcal{N}_+(f)$  has only one compact edge  $\gamma$  and the polynomial  $p(f, \gamma)$  has only one root  $\mu_f$  in  $\mathbb{K}^*$ , i.e.,  $p(f, \gamma) = a_{\gamma}(t - \mu_f)^c$ , and  $a_{\gamma} \in \mathbb{K}^*$ .

## 1.2. On $\nu$ -quasi-ordinary polynomials

Notation 1.4. Let  $f(\boldsymbol{x},z) := \sum c_{\boldsymbol{\alpha},\beta} \boldsymbol{x}^{\boldsymbol{\alpha}} z^{\beta} \in \mathbb{K}[[\boldsymbol{x}]][z], \ \boldsymbol{\alpha} \in \mathbb{N}^d, \ \beta \in \mathbb{N}, \ \text{be a }z$ -polynomial with coefficients in the formal power series ring  $\mathbb{K}[[\boldsymbol{x}]], \ \boldsymbol{x} := (x_1, \ldots, x_d)$ . Since the ring  $\mathbb{K}[[\boldsymbol{x}]][z]$  is factorial, we may assume that  $f(\boldsymbol{x},z) = x_1^{n_1} \cdot \ldots \cdot x_d^{n_d} g(\boldsymbol{x},z)$ , where  $g(\boldsymbol{x},z)$  is regular of order say  $n \geq 0$ , that is,  $g(\boldsymbol{0},z) = a_0 z^n + a_1 z^{n+1} + \ldots, \ a_0 \in \mathbb{K}^*$ . Applying Weierstrass preparation theorem to  $g(\boldsymbol{x},z)$  there exists a unit  $u(\boldsymbol{x},z) \in \mathbb{K}[[\boldsymbol{x},z]]$  and a Weierstrass polynomial  $h(\boldsymbol{x},z) = z^n + a_1(\boldsymbol{x})z^{n-1} + \ldots + a_{n-1}(\boldsymbol{x})z + a_0(\boldsymbol{x})$  with  $a_i(\boldsymbol{0}) = 0$  such that  $g(\boldsymbol{x},z) = h(\boldsymbol{x},z)u(\boldsymbol{x},z)$  and  $u(\boldsymbol{0},0) = a_0 \in \mathbb{K}^*$ , see for example [18, Chapter I, p. 11]. Note that  $\mathcal{N}(g) = \mathcal{N}(h)$  and we do not need to know explicitly h for the constructions in this paper.

Remark 1.5. Let  $f(\boldsymbol{x}, z) := \sum c_{\boldsymbol{\alpha}, \beta} \boldsymbol{x}^{\boldsymbol{\alpha}} z^{\beta} = x_1^{n_1} \cdot \ldots \cdot x_d^{n_d} g(\boldsymbol{x}, z) \in \mathbb{K}[[\boldsymbol{x}]][z]$  be a polynomial and assume we are as in Notation 1.4. Its Newton polyhedron  $\mathcal{N}(f)$  consists only of one vertex  $(\boldsymbol{n}, n) := (n_1, \ldots, n_d, n)$  if and only if  $\mathcal{N}_+(f) = (\boldsymbol{n}, n) + (\mathbb{R}_+)^{d+1}$ .

**Definition 1.6.** A compact face  $\gamma$  of the Newton polyhedron  $\mathcal{N}(f)$  can be eliminated if the polynomial  $f_{\gamma}$  associated with  $\gamma$  can be written as  $f_{\gamma} = \boldsymbol{x}^{\boldsymbol{m}}(z - h(\boldsymbol{x}))^n \in \mathbb{K}[\boldsymbol{x}, z]$ , with  $h \in \mathbb{K}[\boldsymbol{x}]$ ,  $n \ge 1$ . In such a case and by applying the change of variables map  $\sigma \colon \mathbb{K}[[\boldsymbol{x}]][z] \to \mathbb{K}[[\boldsymbol{x}]][z_1]$  defined by  $z \mapsto z_1 = z + h(\boldsymbol{x})$  to f then the face  $\gamma$  is eliminated in the new Newton diagram  $\mathcal{N}_+(f \circ \sigma)$ .

**Definition 1.7.** Let  $f(x, z) = \sum c_{\alpha,\beta} x^{\alpha} z^{\beta} = x_1^{n_1} \cdot \ldots \cdot x_d^{n_d} g(x, z) \in \mathbb{K}[[x]][z]$  be a polynomial and assume we are as in Notation 1.4 and denote  $A := (n, n) \in \mathbb{Q}^{d+1}$ . Let

$$\pi_A \colon \mathcal{N}_+(f) \setminus A \to \mathbb{Q}^d \equiv \mathbb{Q}^{d+1} \cap \{z = 0\}$$

be the projection into  $\mathbb{Q}^d$  with centre at A. We define  $\mathcal{N}_{< n}(f)$  to be the set of points in  $\mathcal{N}_+(f)$  whose z-coordinate is smaller than n. A compact face  $\Gamma$  of  $\mathcal{N}_+(f)$  will be called  $\nu$ -proper if  $A \in \Gamma$ ,  $A \neq \Gamma$  and  $\Gamma \setminus \{A\} \subset \mathcal{N}_{< n}(f)$ . We define  $\mathcal{N}^0(f)$  to be the set of all compact faces of  $\pi_A(\mathcal{N}_{< n}(f))$ .

Remark 1.8. For a  $\nu$ -proper face  $\Gamma$  of  $\mathcal{N}(f)$ ,  $\pi_A(\Gamma)$  is a compact face of  $\mathcal{N}^0(f)$  of dimension one less than the dimension of  $\Gamma$ . In particular,  $\mathcal{N}^0(f)$  consists of exactly one vertex if and only if  $\mathcal{N}(f)$  has only one  $\nu$ -proper compact face and this face is 1-dimensional.

**Definition 1.9.** A regular system of parameters  $x_1, \ldots, x_d, z$  of the local regular ring  $\mathbb{K}[[x, z]]$  is called a *suitable system of coordinates* for a polynomial  $f \in \mathbb{K}[[x]][z]$  if either

- $\mathcal{N}^0(f)$  is void, or
- $\mathcal{N}^0(f)$  has more than one vertex, or
- The set  $\mathcal{N}^0(f)$  consists of exactly one vertex and the corresponding  $\nu$ -proper 1-dimensional compact face  $\Gamma$  cannot be eliminated (Definition 1.6).

Remark 1.10. The condition to be a suitable system of coordinates for a polynomial  $f \in \mathbb{K}[[x]][z]$  is weaker than the conditions of good coordinates in [4] because it involves exactly one  $\nu$ -proper 1-dimensional compact face  $\Gamma$  of the Newton diagram  $\mathcal{N}_+(f)$ , see also Section 3.

**Definition 1.11.** Let  $f(x, z) \in \mathbb{K}[[x]][z]$  be a polynomial and assume we are as in Notation 1.4 and Definition 1.7. We say that f is a  $\nu$ -quasi-ordinary polynomial if

- $\mathcal{N}(f)$  has only one  $\nu$ -proper face  $\Gamma_1$ ,
- $\Gamma_1$  is a 1-dimensional face and
- $\Gamma_1$  cannot be eliminated.

The polynomial  $f_{\Gamma_1}$  is called the *initial form* of f.

The class of  $\nu$ -quasi-ordinary singularities was introduced by H. Hironaka in [19], where he proved that quasi-ordinary polynomials are indeed  $\nu$ -quasi-ordinary polynomials. In comparison with the condition to be a quasi-ordinary polynomial,  $\nu$ -quasi-ordinary is a very mild condition. However,  $\nu$ -quasi-ordinary have interesting properties and some of them will be discussed here.

We consider f to be a  $\nu$ -quasi-ordinary polynomial with its  $\nu$ -proper face  $\Gamma_1$  and keep the above notations. The initial form  $f_{\Gamma_1}$  of f can be written as follows:

$$f_{\Gamma_1} = a_{\Gamma_1} x^n z^{n_{d+1}} \prod_{j=1}^k (z^p - \mu_j x^q)^{m_j},$$
 (1.3)

where factors are irreducible in  $\mathbb{K}[[x]][z]$ , i.e.,  $\gcd(q, p) := \gcd(q_1, \ldots, q_d, p) = 1$ ,  $\mu_j \in \mathbb{K}^*$  with  $\mu_j \neq \mu_i$  if  $i \neq j$ , and  $a_{\Gamma_1} \in \mathbb{K}^*$  (see Proposition 1.2).

**Definition 1.12.** A polynomial  $f(x, z) \in \mathbb{K}[[x]][z]$  is called *elementary* if its Newton polyhedron  $\mathcal{N}(f)$  consists of only one compact face which is a line segment  $[(0, n), (r, 0)] \subset \mathbb{R}^{d+1}$ .

More generally, if  $\Gamma$  is a compact one-dimensional face of the Newton diagram  $\mathcal{N}_+(f)$  of a polynomial  $f(\boldsymbol{x},z) \in \mathbb{K}[[\boldsymbol{x}]][z]$ , we say that the polynomial g is  $\Gamma$ -elementary if its Newton polyhedron  $\mathcal{N}(g)$  consists of one compact face  $\widetilde{\Gamma}$  which is a translation of the face  $\Gamma$ . We denote the initial form of g by  $\operatorname{In}(g) := g_{\widetilde{\Gamma}} \in \mathbb{K}[\boldsymbol{x},z]$ .

In [5, Theorem 1.5] the following factorization theorem was proved.

**Factorization Theorem 1.13.** Let  $f \in \mathbb{K}[[x]][z]$  be a  $\nu$ -quasi-ordinary polynomial with its  $\nu$ -proper face  $\Gamma_1$  such that  $f_{\Gamma_1}$  is as in (1.3). Then there exist k different  $\Gamma_1$ -elementary polynomials  $g_{\Gamma_1,j} \in \mathbb{K}[[x]][z]$ ,  $1 \leq j \leq k$ , which divide f and such that  $\operatorname{In}(g_{\Gamma_1,j}) = (z^p - \mu_j x^q)^{m_j}$ , that is,

$$f = hG^{\Gamma_1}, \quad with \ G^{\Gamma_1} := \prod_{j=1}^k g_{\Gamma_1,j} \ for \ some \ h \in \mathbb{K}[[\boldsymbol{x}]][z].$$
 (1.4)

Factorization Theorem 1.13 has the following consequences.

**Corollary 1.14.** Let f be an irreducible polynomial which is  $\nu$ -quasi-ordinary. Then it is elementary and  $\operatorname{In}(f) = f_0^k$ , where  $f_0 \in \mathbb{K}[x, z]$  is irreducible.

Corollary 1.15. Let f be an elementary polynomial such that  $\operatorname{In}(f) = \psi_1 \cdot \ldots \cdot \psi_r$  is a factorization of  $\operatorname{In}(f)$  where the factors are pairwise coprime. Then we can decompose  $f = f_1 \cdot \ldots \cdot f_r$ , where, for all  $i = 1, \ldots, r$ , the series  $f_i$  is an elementary polynomial such that  $\operatorname{In}(f_i) = \psi_i$ .

Corollary 1.16. Let f be a  $\nu$ -quasi-ordinary polynomial. Then  $f = x^a z^b f^{\Gamma_1} g$ , where  $f^{\Gamma_1}$  is  $\Gamma_1$ -elementary, and  $\operatorname{In}(f^{\Gamma_1}) = f_{\Gamma_1}$ .

**1.3.** On Newton maps and the Newton process. We need to introduce more notations in order to define *Newton maps*. For every  $1 \le i \le d$ , let

$$c_i := \gcd(p, q_i), \quad p_i := \frac{p}{c_i}, \quad \text{and} \quad q'_i := \frac{q_i}{c_i}.$$
 (1.5)

We will consider the following three maps. The first map depends both on the face  $\Gamma_1$  and on the root  $\mu_j$ , see (1.3). Let  $(\boldsymbol{u}, u) \in \mathbb{N}^{d+1}$  be integers such that  $1 + \boldsymbol{u} \cdot \boldsymbol{q} = up$ . Let  $\delta_{\Gamma_1,j}$  be the map

$$\begin{split} \delta_{\Gamma_1,j} \colon \mathbb{K}[[x]][z] &\to \mathbb{K}[[x]][z] \\ h(x,z) &\mapsto h(\mu_j^{u_1} x_1, \, \dots, \, \mu_j^{u_d} x_d, \, z). \end{split}$$

Then

$$\delta_{\Gamma_1,j}(z^p - \mu_j \boldsymbol{x}^{\boldsymbol{q}}) = z^p - \mu_j^{up} \boldsymbol{x}^{\boldsymbol{q}}.$$

The second map  $\epsilon_{\Gamma_1}$  depends only on the face  $\Gamma_1$ :

$$\epsilon_{\Gamma_1} \colon \mathbb{K}[[\boldsymbol{x}]][z] \to \mathbb{K}[[\boldsymbol{y}]][z_1]$$

$$h(\boldsymbol{x}, z) \mapsto h(y_1^{p_1}, \dots, y_d^{p_d}, \boldsymbol{y}^{\boldsymbol{q}'}z_1).$$

It is easily seen that

$$(\epsilon_{\Gamma_1} \circ \delta_{\Gamma_1,j})(z^p - \mu_j \boldsymbol{x}^{\boldsymbol{q}}) = \boldsymbol{y}^{p\boldsymbol{q}'}(z_1^p - \mu_j^{up}).$$

Now the third map  $\tau_j$  will depend on the root  $\mu_j$ :

$$\tau_j \colon \mathbb{K}[[\boldsymbol{y}]][z_1] \to \mathbb{K}[[\boldsymbol{y}]][z_2]$$

$$h(\boldsymbol{y}, z_1) \mapsto h(\boldsymbol{y}, z_2 + \mu_i^u).$$

Note that:

$$(\tau_j \circ \epsilon_{\Gamma_1} \circ \delta_{\Gamma_1,j})(z^p - \mu_j \boldsymbol{x}^{\boldsymbol{q}}) = \boldsymbol{y}^{\boldsymbol{p} * \boldsymbol{q}}((z_2 + \mu_i^u)^p - \mu_i^{up})$$

and the last factor is of order one in  $z_2$ . Note that  $p_iq_i = pq'_i$  for  $1 \le i \le d$ , i.e, p \* q = pq'.

**Definition 1.17.** The Newton map  $\sigma_{\Gamma_1,j}$  associated with  $\Gamma_1$  and the root  $\mu_j$  is the composition map  $\sigma_{\Gamma_1,j} := \tau_j \circ \epsilon_{\Gamma_1} \circ \delta_{\Gamma_1,j} : \mathbb{K}[[\boldsymbol{x}]][z] \to \mathbb{K}[[\boldsymbol{y}]][z_2]$ .

Also the following statement was proved in [5, Lemma 2.2].

**Lemma 1.18.** After the Newton map  $\sigma_{\Gamma_1,j}$ , the total transform  $f_{\Gamma_1,j} := \sigma_{\Gamma_1,j}(f) \in \mathbb{K}[[y]][z_2]$  of the polynomial f can be written as

$$f_{\Gamma_1,j}(\boldsymbol{y},z_2) = \boldsymbol{y}^{m\boldsymbol{p}*\boldsymbol{q}+\boldsymbol{n}*\boldsymbol{p}} f_{1,\Gamma_1,j}(\boldsymbol{y},z_2),$$

where  $f_{1,\Gamma_1,j}(\mathbf{0}, z)$  is regular of order  $m_j$ , and  $m := \sum_{j=1}^k m_j$ . Moreover by chain rule

$$\frac{\partial (f \circ \sigma_{\Gamma_1,j})}{\partial z_2}(\boldsymbol{y},\,z_2) = \boldsymbol{y}^{\boldsymbol{q}'}\left(\frac{\partial f}{\partial z} \circ \sigma_{\Gamma_1,j}\right)(\boldsymbol{y},\,z_2).$$

- 1.3.1. Newton's process associated with the  $\nu$ -proper face  $\Gamma_1$ . We can perform a change of variables of the type  $z \mapsto z + h(x)$  in order to have a suitable system of coordinates for  $f_{1,\Gamma_1,j}$ ; if  $f_{1,\Gamma_1,j}$  is again  $\nu$ -quasi-ordinary one can iterate the process until one gets either a monomial times a unit or a not  $\nu$ -quasi-ordinary polynomial. This process is called Newton's process; note that  $m_j < n$  because the face  $\Gamma_1$  cannot be eliminated.
- 1.3.2. Factorization and the Newton process associated with other compact 1-dimensional faces of the Newton polyhedron. Since f is a  $\nu$ -quasi-ordinary polynomial with  $\nu$ -proper face  $\Gamma_1$  then  $\Gamma_1$  is a 1-dimensional face of its Newton diagram  $\mathcal{N}_+(f)$ . We assume that  $\Gamma_1$  is the segment  $[A, A_1]$ . If the z-coordinate of  $A_1$  is  $n^1 > 0$ , we denote by  $\mathcal{N}_{< n^1}(f)$  the set of points in  $\mathcal{N}_+(f)$  whose z-coordinate is smaller than  $n^1$ . Let  $\pi_{A_1} : \mathcal{N}_{< n^1}(f) \setminus A_1 \to \mathbb{Q}^d$  be the projection into with center  $A_1$  and let  $\mathcal{N}^{0,1}(f)$  be the convex hull of the image by  $\pi_{A_1}$  of  $\mathcal{N}_{< n^1}(f)$ . If  $\mathcal{N}^{0,1}(f)$  has only one vertex then there is another face  $\Gamma_2$  of the Newton diagram  $\mathcal{N}_+(f)$  which is of dimension 1. We go further on this construction with  $\Gamma_1 = [A, A_1], \Gamma_2 = [A_1, A_2], \ldots, \Gamma_s = [A_s, A_{s+1}]$  until one of the following cases arises:

(NW1) The z-coordinate  $n^{s+1}$  of  $A_{s+1}$  is zero.

(NW2)  $\mathcal{N}^{0,s+1}(f) = \varnothing$ .

(NW3)  $\mathcal{N}^{0,s+1}(f)$  has more than one vertex.

Moreover  $\Gamma_1 \cup \Gamma_2 \cup \cdots \cup \Gamma_s$  is a monotone polygonal path in  $\mathcal{N}(f)$ .

**Lemma 1.19.** The Newton polyhedron of f is not a monotone polygonal path if and only if  $\mathcal{N}^{0,s+1}(f)$  has more than one vertex, i.e., (NW3) arises.

For every edge  $\Gamma_{\ell}$  of  $\Gamma_2, \ldots, \Gamma_s$ , the initial form  $f_{\Gamma_{\ell}}$  of f can be written as in (1.3) as follows:

$$f_{\Gamma_{\ell}} = a_{\Gamma_{\ell}} \boldsymbol{x}^{n^{\ell}} z^{n_{d+1}^{\ell}} \prod_{j=1}^{k^{\ell}} (z^{p_{\ell}} - \mu_{j}^{\ell} \boldsymbol{x}^{q^{\ell}})^{m_{j}^{\ell}},$$
(1.6)

the factors being irreducible in  $\mathbb{K}[[x]][z]$ , i.e.,  $\gcd(q^\ell, p_\ell) = 1$ ,  $\mu_j^\ell \in \mathbb{K}^*$  with  $\mu_j^\ell \neq \mu_i^\ell$  and  $a_{\Gamma_\ell} \in \mathbb{K}^*$ . For each root  $\mu_j^\ell$  of its face polynomial  $f_{\Gamma_\ell}$  one applies the corresponding Newton map  $\sigma_{\Gamma_\ell,\mu_j^\ell}$ . At each step, we encode the information given by the corresponding Newton diagram. The process stops because the z-degree decreases since we are in a suitable system of coordinates. In next section Newton trees of a polynomial  $f(x, z) \in \mathbb{K}[[x]][z]$  are constructed by recursion on the number of steps of Sections 1.3.1 and 1.3.2.

One can also apply recursively Factorization Theorem 1.13 to h, see (1.4), to get  $k^{\ell}$  different  $\Gamma_{\ell}$ -elementary polynomials  $g_{\Gamma_{\ell},j} \in \mathbb{K}[[x]][z]$ ,  $1 \leq j \leq k^{\ell}$ , such that  $G^{\Gamma_{\ell}} := \prod_{j=1}^{k^{\ell}} g_{\Gamma_{\ell},j}$  divides f, that is,

$$f = H(\boldsymbol{x}, z) \prod_{t=1}^{s} G^{\Gamma_t}$$
 for some  $H \in \mathbb{K}[[\boldsymbol{x}]][z]$ . (1.7)

#### 2. The Newton Tree of a Polynomial

**2.1.** Construction of the Newton tree. Let  $f(x, z) \in \mathbb{K}[[x]][z]$  be a polynomial and we assume that we are as in Notation 1.4. In this section we associate with f(x, z) a tree  $\mathcal{T}(f)$  called Newton tree of f because its first steps are built using both its Newton diagram  $\mathcal{N}_+(f)$  and the set  $\mathcal{N}^0(f)$  of compact faces of  $\pi(\mathcal{N}_{< n}(f))$ . Further steps will be based on the Newton process associated with f (see subsections 1.3.1 and 1.3.2).

For a given f, we are going to associate a tree  $\mathcal{T}_{\mathcal{N}}(f)$ , called *vertical tree*. First, if f is not in a suitable system of coordinates (see Definition 1.9), we perform a change of variables such that it is the case; in order to simplify the notations we denote again by f the resulting polynomial. We keep the notations of Definition 1.7. The tree  $\mathcal{T}_{\mathcal{N}}(f)$  is built using its Newton diagram  $\mathcal{N}_{+}(f)$ .

We distinguish three cases.

Case 1. The Newton polyhedron  $\mathcal{N}(f)$  consists only in one vertex (see Remark 1.5). Then the Newton tree  $\mathcal{T}_{\mathcal{N}}(f)$  of f is given in Figure 1(a).

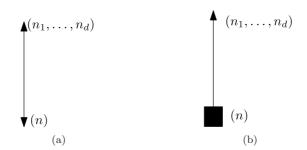


FIGURE 1.

Case 2. The set  $\mathcal{N}^0(f)$  of all compact faces of  $\pi(\mathcal{N}_{< n}(f))$  has more than one vertex. Then the Newton tree  $\mathcal{T}_{\mathcal{N}}(f)$  of f is given in Figure 1(b).

Case 3. The set  $\mathcal{N}^0(f)$  has exactly one vertex. Since we are in a suitable system of coordinates for f, the face  $\Gamma_1$  cannot be eliminated, (i.e., f is a  $\nu$ -quasi-ordinary polynomial). In such a case we are as in subsection 1.3.2, and there exist s compact edges  $\Gamma_1 := [A, A_1], \Gamma_2 := [A_1, A_2], \ldots, \Gamma_s := [A_s, A_{s+1}]$  of the Newton polyhedron  $\mathcal{N}(f)$  until one of the cases (NW1), (NW2) or (NW3) happen.

Furthermore  $\Gamma_1 \cup \Gamma_2 \cup \cdots \cup \Gamma_s$  is a monotone polygonal path in  $\mathcal{N}(f)$ . With this monotone polygonal path we associate a decorated vertical graph  $\mathcal{T}_{\mathcal{N}}(f)$  (which depends on  $\mathcal{N}(f)$  in a suitable system of coordinates for f). With each compact 1-dimensional face  $\Gamma_\ell$  of the polygonal path, we associate a vertex  $v_\ell$ . If two compact faces intersect at one point we draw a vertical edge from one vertex to the other. Thus these vertices are drawn on a vertical line by the increasing order of the slopes, i.e.,  $v_1, v_2, \ldots, v_s$  from above to below in order. Decorations of this vertical line  $\mathcal{T}_{\mathcal{N}}(f)$  are as follows:

- On the top of the vertical line we add an arrow-head decorated with  $n := (n_1, \ldots, n_d)$ .
- On the bottom of the vertical line, we add:
  - an arrow-head decorated with  $(n^{s+1})$  in cases (NW1)  $(n^{s+1} = 0)$  and (NW2);
  - a black box decorated with  $(n^{s+1})$  in case (NW3).
- for an edge  $\Gamma_{\ell}$  which is defined in coordinates  $(\boldsymbol{\alpha}, \beta) := (\alpha_1, \ldots, \alpha_d, \beta)$  of  $\mathbb{R}^{d+1}$  by the intersection of d hyperplanes of equations

$$p_{\ell}\alpha_k + q_k^{\ell}\beta = N_k^{\ell}, \quad 1 \leqslant k \leqslant d, \quad \gcd(\boldsymbol{q}^{\ell}, p_{\ell}) = 1, \tag{2.1}$$

the corresponding vertex  $v_\ell$  supports the following decorations:

- The vertex itself is decorated with  $((N_1^{\ell}, \ldots, N_d^{\ell}))$ .
- The lower edge is decorated near  $v_{\ell}$  with  $p_{\ell}$ .
- The upper edge is decorated near  $v_{\ell}$  with  $(q_1^{\ell}, \ldots, q_d^{\ell})$ .

We describe now the construction of  $\mathcal{T}(f)$ . Recall that we assume that f is in suitable coordinates.

- **Step 1.** If f is in either Case 1 or 2 then  $\mathcal{T}(f) := \mathcal{T}_{\mathcal{N}}(f)$ . If f is in Case 3 we continue the process.
- **Step 2.** For every 1-dimensional face  $f_{\Gamma_{\ell}}$  and for each root  $\mu_j^{\ell}$  of its face polynomial  $f_{\Gamma_{\ell}}$  one applies the corresponding Newton map  $\sigma_{\Gamma_{\ell},\mu_j^{\ell}}$  to f and get a polynomial  $f_{\Gamma_{\ell},j}$ ; we perform a change of coordinates to be in suitable coordinates.
- Step 3. The tree  $\mathcal{T}(f)$  will be obtained by gluing in a suitable way the tree  $\mathcal{T}_{\mathcal{N}}(f)$  and  $\mathcal{T}(f_{\Gamma_{\ell},j})$ , which may be assumed constructed recursively. The tree  $\mathcal{T}(f_{\Gamma_{\ell},j})$  will be attached to the vertex  $v_{\ell}$  of  $\mathcal{T}_{\mathcal{N}}(f)$  by a horizontal edge which links  $v_{\ell}$  to the top vertex of  $\mathcal{T}_{\mathcal{N}}(f_{\Gamma_{\ell},j})$ , where the top arrow has been deleted.

Now to decorate the tree  $\mathcal{T}(f)$  we need more definitions. Let v be a vertex on  $\mathcal{T}(f)$ . If v is on the first left vertical line, we say that v has no preceding vertex. If v is not on the first left vertical line, the vertical line on which v lies, is linked by an horizontal edge ending on a vertex  $v_1$ , to a vertical line. Then  $v_1$  is said to be the preceding vertex of v. Note that the path between  $v_1$  and v can have many vertical edges, but has exactly one horizontal edge. Now denote by  $\mathcal{S}(v) = \{v_i, v_{i-1}, \ldots, v_1, v_0 = v\}$ , where  $v_j$  is the preceding vertex of  $v_{j-1}$  for  $j = 1, \ldots, i$  and  $v_i$  has no preceding vertex. We say that  $\mathcal{S}(v) \setminus \{v\}$  is the set of preceding vertices of v. Now in the contruction of the Newton tree  $\mathcal{T}(f)$ , we glue on  $\mathcal{T}_{\mathcal{N}}(f)$  at  $v_i$  a tree  $\mathcal{T}(f_{\Gamma_i,l})$ , where  $\mathcal{S}(v) = \{v_{i-1}, \ldots, v_1, v_0 = v\}$ . Assume the decorations of the edges attached to  $v_i$  on  $\mathcal{T}_{\mathcal{N}}(f)$  are  $((q_1^i, \ldots, q_d^i), p_i)$ .

We are going to decorate the edges attached to  $v_0$  in  $\mathcal{T}(f)$  with  $((Q_1^0, \ldots, Q_d^0), p_0)$  following these rules:

• The decoration  $p_0$  coincides with the corresponding decoration in the vertical tree containing  $v_0$ .

• Let us assume that the decorations of  $v_j$  on  $\mathcal{T}(f_{\Gamma_i,l})$  are  $((Q_1^{j,i},\ldots,Q_d^{j,i}),p_j)$  for  $j=0,\ldots,i-1$ . Then

$$Q_k^0 := \frac{p_i q_k^i p_{i-1}^2 \cdot \dots \cdot p_1^2 p_0}{\gcd(p_i, q_k^i) \cdot \gcd(p_{i-1}, Q_k^{i-1,i}) \cdot \dots \cdot \gcd(p_1, Q_k^{1,i})} + Q_k^{0,i},$$

$$k \in \{1, \dots, d\}.$$
(2.2)

Note that, in particular,  $Q_k^i = q_k^i$ .

It is useful to have another way for computing these decorations, see [5, (5.1)]. Let us assume with the above notations that the decorations  $((Q_1^j, \ldots, Q_d^j), p_j)$  for  $\mathcal{T}(f)$  have been defined for  $v_j$ ,  $j=0,1,\ldots,i-1$ . Recall that  $Q_k^i=q_k^i$ . Let us denote  $((q_1^0,\ldots,q_d^0),p_0)$  the decorations of  $v_0$  in its vertical tree.

Lemma 2.1. With the above notations

$$Q_k^0 = \frac{p_1 Q_k^1 p_0}{\gcd(p_1, Q_k^1)} + q_k^0.$$

*Proof.* Let us assume first that i=1. In this case  $Q_k^{0,1}=q_k^0$  since  $\mathcal{T}(f_{\Gamma_1,l})=\mathcal{T}_{\mathcal{N}}(f_{\Gamma_1,l})$  and  $Q_k^1=q_k^1$  since i=1. The result follows from a direct substitution in (2.2).

If i > 1 then, by induction hypothesis, we have  $Q_k^{0,i} = \frac{p_1 Q_k^{1,i} p_0}{\gcd(p_1, Q_k^{1,i})} + q_k^0$ . Hence,

$$\begin{split} Q_k^0 &= \frac{p_i q_k^i p_{i-1}^2 \cdot \ldots \cdot p_1^2 p_0}{\gcd(p_i, q_k^i) \cdot \gcd(p_{i-1}, Q_k^{i-1,i}) \cdot \ldots \cdot \gcd(p_1, Q_k^{1,i})} + Q_k^{0,i} \\ &= \frac{p_i q_k^i p_{i-1}^2 \cdot \ldots \cdot p_1^2 p_0}{\gcd(p_i, q_k^i) \cdot \gcd(p_{i-1}, Q_k^{i-1,i}) \cdot \ldots \cdot \gcd(p_1, Q_k^{1,i})} + \frac{p_1 Q_k^{1,i} p_0}{\gcd(p_1, Q_k^{1,i})} + q_k^0 \\ &= \Big(\frac{p_i q_k^i p_{i-1}^2 \cdot \ldots \cdot p_2^2 p_1}{\gcd(p_i, q_k^i) \cdot \gcd(p_{i-1}, Q_k^{i-1,i}) \cdot \ldots \cdot \gcd(p_1, Q_k^{1,i})} + Q_k^{1,i}\Big) \frac{p_1 p_0}{\gcd(p_1, Q_k^{1,i})} + q_k^0. \end{split}$$

Using (2.2)

$$Q_k^1 = \frac{p_i q_k^i p_{i-1}^2 \cdot \dots \cdot p_2^2 p_1}{\gcd(p_i, q_k^i) \cdot \gcd(p_{i-1}, Q_k^{i-1,i}) \cdot \dots \cdot \gcd(p_1, Q_k^{1,i})} + Q_k^{1,i}$$

and the result follows.

Remark 2.2. By construction the Newton tree  $\mathcal{T}(f)$  has vertical parts and horizontal parts: vertical parts correspond to Newton diagrams of total transforms by Newton maps and horizontal edges are edges used for connecting vertical parts.

**Definition 2.3.** An *end* of the Newton tree  $\mathcal{T}(f)$  is either an arrow-head or a black box. The arrow-heads or black boxes decorated with (0) will be called *dead ends*.

Remark 2.4. By construction black boxes in the Newton tree  $\mathcal{T}(f)$  appear if and only if we reach at some step Case 2. This means that the Newton tree  $\mathcal{T}(f)$  ends with arrow-heads if and only if we never reach Case 2.

**Definition 2.5.** For every vertex  $v_j$  of the Newton tree  $\mathcal{T}(f)$ , the global numerical data of the vertex  $v_j$  is  $(\mathbf{N}_{v_j}, p_j)$ , where  $\mathbf{N}_{v_j} := (N_1^j, \dots, N_d^j)$ .

Remark 2.6. Global numerical data of the first vertex satisfy the following property. If  $\sigma$  is a Newton map associated with this vertex, then

$$\sigma(f) = \mathbf{y}^{\frac{N}{c}} f_1(\mathbf{y}, z_1),$$

where no  $y_i$  divides  $f_1$  and  $c_i = \gcd(p, q_i)$  as in (1.5).

**Definition 2.7.** The numerical data  $((Q_1^j, \ldots, Q_d^j), p_j)$  associated with each vertex  $v_j$  of the Newton tree  $\mathcal{T}(f)$  are called the *local numerical data* of the vertex  $v_j$ . The *gcd of the vertex*  $v_j$  is  $\mathbf{c} := (\gcd(Q_1^j, p_j), \ldots, \gcd(Q_d^j, p_j))$ .

Remark 2.8. It is easy to show that, if  $((Q_1, \ldots, Q_d), p)$  are the local numerical data of a vertex v and  $((Q'_1, \ldots, Q'_d), p_j)$  are the local numerical data of a vertex v' such that v is the preceding vertex of v'. Then, applying Lemma 2.1, the following strict inequalities hold:

$$Q'_i > \frac{pQ_i}{\gcd(p, Q_i)} p_j$$
 for  $i = 1, ..., d$  and  $j = 1, ..., r$ .

This condition is called the growth condition on the local numerical data.

Remark 2.9. In the case for d = 1, and if we forget about vertical and horizontal edges the Newton trees decorated with local numerical data are the Eisenbud-Neumann diagrams of the corresponding germ defined in [11].

**Definition 2.10.** The *valency* of a vertex in a Newton tree is the number of edges attached to the vertex.

**2.2.** Comparable polynomials and colored Newton trees. Suppose that  $f(x, z), g(x, z) \in \mathbb{K}[[x]][z]$  are polynomials. We assume that we are as in Notation 1.4 and Definition 1.7.

**Definition 2.11.** Two polynomials f and g in  $\mathbb{K}[[x]][z]$  are called *comparable* if their resultant  $\operatorname{Res}_z(f, g)(x)$  of f and g with respect to z is equal to a monomial times a unit, that is,

$$\operatorname{Res}_{z}(f, g)(\boldsymbol{x}) = \boldsymbol{x}^{\boldsymbol{n}}u(\boldsymbol{x}), \quad \text{with } \boldsymbol{n} \in \mathbb{N}^{d}, u(\boldsymbol{x}) \in \mathbb{K}[[\boldsymbol{x}]], \text{ and } u(\boldsymbol{0}) \neq 0.$$

2.2.1. Colored Newton trees. Colored Newton trees are associated with the product of two polynomials  $f(\mathbf{x}, z)$ ,  $g(\mathbf{x}, z) \in \mathbb{K}[[\mathbf{x}]][z]$ . We assume fg is in suitable coordinates. Take the Newton tree  $\mathcal{T}(fg)$  of the product  $fg = f(\mathbf{x}, z) \cdot g(\mathbf{x}, z)$  and add two colors to it, say red and blue. Blue color is associated with f and red color with g. The colored Newton tree  $\mathcal{T}(fg)$  can have blue parts, red parts and blue-red parts.

**Definition 2.12.** We consider two polynomials  $f(x, z) := x^n f_1(x, z)$ ,  $g(x, z) := x^m g_1(x, z) \in \mathbb{K}[[x]][z]$  such that  $f_1(x, z)$  and  $g_1(x, z)$  are regular of order say  $m, n \ge 0$ . In such a case the product

$$f_1g_1(\mathbf{0}, z) = a_0 z^{m+n} + \dots$$

satisfies  $a_0 \neq 0$ . In this case  $A := (\mathbf{0}, m+n) \in \mathcal{N}_+(f_1g_1)$  and let  $\pi$  be the projection into  $\mathbb{Q}^d$  with centre A as before and we consider  $\mathcal{N}_0(f_1g_1)$  as in Definition 1.7. We consider the Newton diagram  $\mathcal{N}_+(fg) = \mathcal{N}_+(\boldsymbol{x}^{n+m}) + \mathcal{N}_+(f_1g_1)$ , which is used to

construct the first vertical part  $\mathcal{T}_{\mathcal{N}}(fg)$  of the Newton tree  $\mathcal{T}(fg)$  of fg. Three possible cases may arise either

- (1)  $\mathcal{T}_{\mathcal{N}}(fg)$  is blue colored if deg  $g_1(\mathbf{0}, z) = 0$  or,
- (2)  $\mathcal{T}_{\mathcal{N}}(fg)$  is red colored if deg  $f_1(\mathbf{0}, z) = 0$  or,
- (3)  $\mathcal{T}_{\mathcal{N}}(fg)$  is bi-colored blue-red otherwise.

We apply the same rule for every steps of the Newton process. In particular every vertical line in the Newton tree  $\mathcal{T}(fg)$  of fg has the same (bi)color. Bicolored vertices of the bicolored Newton tree  $\mathcal{T}(fg)$  will be called *common vertices* of f and g.

**Example 2.13.** The bi-colored Newton tree  $\mathcal{T}(fg)$  of fg, where  $f = z^2 - x^3$  and  $g = z^3 - x^2$ , is as in Figure 2 where all vertical lines are bi-colored, the above horizontal line is red and the below one is blue.

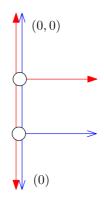


FIGURE 2.

**Definition 2.14.** Two polynomials f and g in  $\mathbb{K}[[x]][z]$  are called *separated* if there exists a suitable system of coordinates for fg such that all ends which are not deadends of the bi-colored Newton tree  $\mathcal{T}(fg)$  are either red colored or blue colored, see Definition 2.3.

Remark 2.15. If  $d=1,\ f$  and g are separated if they do not share a common component.

**Definition 2.16.** Let  $f, g \in \mathbb{K}[[x]][z]$  such that the system of coordinates is suitable for fg. Let v a vertex of the Newton tree  $\mathcal{T}(fg)$  corresponding to an edge  $\Gamma$ . Let  $\hat{f}_{\Gamma}$  obtained from  $f_{\Gamma}$  by taking away the powers of  $x_i$  and z. We say that f and g separate at v if  $\gcd(\hat{f}_{\Gamma}, \hat{g}_{\Gamma})$  is not equal to either  $\hat{f}_{\Gamma}$  or  $\hat{g}_{\Gamma}$ . The order of separation of f at v is the z-degree of  $\frac{\hat{f}_{\Gamma}}{\gcd(\hat{f}_{\Gamma}, \hat{g}_{\Gamma})}$ .

**Example 2.17.** We illustrate the two kinds of separation of f and g. In Figure 3(a) after two Newton maps the total transform of fg, see Lemma 1.18, has a Newton polyhedron with two different compact 1-dimensional faces and each of them corresponds either to (the total transform of) f and the other to (the total

transform of) g. In Figure 3(b) they have different polynomials on the same face, see (1.3).

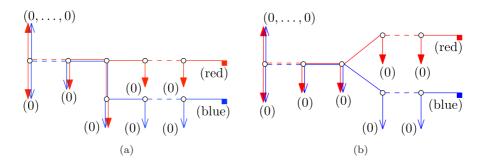


FIGURE 3.

Remark 2.18. If two polynomials  $f, g \in \mathbb{K}[[x]][z]$  are separated then there is at least a vertical line in the bi-colored Newton tree  $\mathcal{T}(fg)$  which corresponds to a Newton polygon of a step of its Newton process, after which its Newton tree  $\mathcal{T}(fg)$  is not bi-colored, see for example Figures 2, 3(a) or 3(b).

Remark 2.19. Let  $f, g \in \mathbb{K}[[x]][z]$  be two polynomials. Assume there exists a suitable system of coordinates of each of them. For a common vertex of the bicolored Newton tree  $\mathcal{T}(fg)$  its local numerical data is the same as its local numerical data as vertex of  $\mathcal{T}(f)$  and the same as its local numerical data as vertex of  $\mathcal{T}(g)$ .

Next theorem is one of the main results in [5].

**Theorem 2.20** [5, Theorem 1.6]. If two polynomials f and g in  $\mathbb{K}[[x]][z]$  are separated then they are comparable, i.e., their resultant  $\operatorname{Res}_z(f, g) = x^{\alpha} \eta(x)$  with  $\eta(x) \in \mathbb{K}[[x]], \alpha \in \mathbb{N}^d$  and  $\eta(0) \neq 0$ .

**2.3.** Computation of resultants. Let f and g be two Weierstrass polynomials which are comparable polynomials and such that their Newton trees have only one end which is not a dead-end. The main result in this section is to show that the resultant  $\operatorname{Res}_z(f,g)$  can be read from the colored Newton tree of fg decorated with its local numerical data. This result can be seen as a generalization of [13, Corollaire 5], where P. González Pérez gave information about the Newton diagram of the resultant of two quasi-ordinary hypersurfaces satisfying an appropriate non-degeneracy condition.

Let f and g be two Weierstrass polynomials in  $\mathbb{K}[[x]][z]$  and asume we are as in Notation 1.4 and the system of coordinates is suitable for both f and g. We consider the colored Newton tree  $\mathcal{T}(fg)$  of fg decorated with its local numerical data. We recall the following result [5], where it was shown that to compute their resultant  $\operatorname{Res}_z(f,g)$  one can use the factorization of total transforms of f and g after Newton process even if they are not factorizations of f and g.

**Proposition 2.21** [5, Proposition 5.11]. Let  $f, g \in \mathbb{K}[[x]][z]$  be two Weierstrass polynomials in suitable coordinates which are comparable and such that their Newton

trees  $\mathcal{T}(f)$  and  $\mathcal{T}(g)$  have only one end which is not a dead-end. Then,  $\operatorname{Res}_z(f, g)$  can be read from the colored Newton tree  $\mathcal{T}(fg)$  of fg decorated with its local numerical data.

More precisely, for every  $1 \leq i \leq d$  the multiplicity of  $x_i$  in  $\operatorname{Res}_z(f, g)$  is computed as follows:

- (1) Consider the path from the blue-end representing f to the red-end representing q.
- (2) Take the product of all the decorations that are adjacent to the path (the *i*-th coordinate in the case of d-uples), including decorations of these ends.
- (3) Multiply it by the the i-th coordinate of the gcd of the vertices before the vertex where they separate, see Definition 2.7.

### 3. P-GOOD COORDINATES

Let  $f(x, z) \in \mathbb{K}[[x]][z]$  be a polynomial and we assume that we are as in Notation 1.4 and in a suitable system of coordinates for f. There is a system of good coordinates, introduced by P. González-Pérez in [14] and [16, Lemma 3.2], that is called P-good coordinates.

**Definition 3.1.** A polynomial  $f(x, z) \in \mathbb{K}[[x]][z]$  is in *P-good coordinates* if

- (1) its Newton polyhedron  $\mathcal{N}(f)$  is a monotone polygonal path,
- (2) if there exists a face  $\Gamma$  of  $\mathcal{N}(f)$  whose line supporting  $\Gamma$  is given by equations  $\{\alpha_i+q_i\beta=N_i, i=1,\ldots d\}$ , then  $\Gamma$  is unique and hits the hyperplane  $\beta=0$ ,
- (3) if there exists a face  $\Gamma$  of  $\mathcal{N}(f)$  whose line supporting  $\Gamma$  is given by equations  $\{\alpha_i + q_i\beta = N_i, i = 1, \dots d\}$ , then the corresponding polynomial  $f_{\Gamma}$  is not of the form  $a_{\Gamma} \boldsymbol{x}^{\boldsymbol{m}} (z h(\boldsymbol{x}))^n$  with  $h(\boldsymbol{x}) \in \mathbb{K}[\boldsymbol{x}]$  and  $a_{\Gamma} \in \mathbb{K}^*$ , see (1.3).

Let us assume that there is a suitable system of coordinates such that the Newton polyhedron of f is a monotone polygonal path. Let  $\mathcal{T}(f)$  be the Newton tree of f in this suitable system of coordinates.

**Proposition 3.2.** Starting from a suitable system of coordinates such that the Newton polyhedron of f is a monotone polygonal path we find a P-good system of coordinates for f by a change of coordinates of the type  $z \mapsto z + ax^q$ . The Newton tree in this system of coordinates can be deduce from  $\mathcal{T}(f)$  in a unique way.

*Proof.* Let  $\mathcal{T}_{\mathcal{N}}(f)$  be the first vertical line of  $\mathcal{T}(f)$ . The vertices are denoted by  $v_1, \ldots, v_n$  from top to bottom decorated with (q, p).

Case 1. None of the decorations p of the vertices of  $\mathcal{T}_{\mathcal{N}}(f)$  is equal to 1

In this case, f is already in P-good coordinates.

Case 2. There is a decoration p of some vertex which is equal to one.

Let i be the smallest index such that  $p_i = 1$ . We have two cases to consider.

Sub-Case 1.  $i \neq n$ .

In this case, f is not in P-good coordinates. Assume the decorations of the vertex  $v_i$  in  $\mathcal{T}_{\mathcal{N}}(f)$  are  $(q_1^i, \ldots, q_d^i, p_i = 1)$  and the face polynomial is

$$\boldsymbol{x}^{k'} z^k \prod_i (z - \mu_j \boldsymbol{x}^{\boldsymbol{q}^i})^{m_j}.$$

Consider the change of coordinates

$$z = z' + ax^{q^i},$$

where  $a \neq \mu_j$  for all j. Under this change of variables neither the faces  $\Gamma_1, \ldots, \Gamma_{i-1}$  nor their face polynomials do change. The face  $\Gamma_i$  has the same equation but now the face polynomial is

$$x^{k'}(z' + ax^{q^i}) \prod_{i} (z - (\mu_j - a)x^{q^i})^{m_j}.$$

The new face  $\Gamma'_i$  hits the *x*-hyperplane. Then f is in P-good coordinates.

Now we compare the tree  $\mathcal{T}(f)$  with the tree  $\mathcal{T}'(f)$  in these new coordinates. The Newton maps corresponding to the vertices  $v_1, \ldots, v_{i-1}$  are the same.

At  $v_i$ , we have to consider the Newton maps

$$x = y$$
,  $z = y^{q^i}(z_1 - \mu_j)$ .

The Newton maps to be considered at  $v'_i$  are given by the equations

$$x = y, \quad z' = y^{q^i}(z'_1 - (\mu_j - a)).$$

Then  $z_1 = z'_1$  and along the corresponding edge nothing is changed. Now at  $v'_i$  we have also to consider the Newton map

$$\boldsymbol{x} = \boldsymbol{y}, \ z' = \boldsymbol{y}^{\boldsymbol{q}^i}(z_1' - a) \quad \Rightarrow \quad \boldsymbol{x} = \boldsymbol{y}, \ z = \boldsymbol{y}^{\boldsymbol{q}^i}z_1'.$$

In the change of variables, the monomial  $x^{\alpha}z^{\beta}$  becomes  $y^{\alpha+\beta q^i}z_1^{\prime\beta}$ .

We consider the transformation in the affine space

$$(\boldsymbol{\alpha}, \beta) \mapsto (\boldsymbol{\alpha} + \beta \boldsymbol{q}^i, \beta).$$

A hyperplane with equations  $\{P\alpha_j + Q_j\beta = 0\}$  maps to  $\{P\alpha'_j + (Q_j - Pq_j^i)\beta = 0\}$ . Then the faces  $\Gamma_{i+1}, \ldots, \Gamma_n$ , map to faces of the Newton diagram in the coordinates  $(\boldsymbol{y}, z'_1)$  and since  $\Gamma_{i+1}, \ldots, \Gamma_n$  was a monotone polygonal path, it transforms to a monotone polygonal path.

Now consider a face polynomial  $x^{k'}z^k\prod_l(z^P-\mu_lx^Q)^{m_l}$ . In the change of coordinates it becomes

$$\boldsymbol{y}^{k'+\boldsymbol{q}^ik"}z_1'^k\prod_l(z^P-\mu_l\boldsymbol{y}^{\boldsymbol{Q}-\boldsymbol{q}^iP})^{m_l}.$$

We consider the following Newton maps. Let (u, u) be integers such that  $1+u\cdot Q=uP$ . We have  $1+u\cdot (Q-Pq^i)=(u-u\cdot q^i)P$ . The Newton maps are

$$x = \mu_l^u x'^P$$
,  $z = x'^Q (z_2 + \mu_j^u)$  and  $y = \mu_l^u y'^P$ ,  $z'_1 = y'^{(Q-Pq^i)} (z'_2 + \mu_j^{(u-u \cdot q^i)})$ .

Since  $z = x^{q^i} z_1'$ , we have x' = y',  $z_2 = \mu_l^{u \cdot q^i} z_2'$ . Then we have essentially the same coordinates.

In conclusion, to get the tree of f in P-good coordinates from  $\mathcal{T}(f)$ , we have to cut the edge  $e_i$  under  $v_i$ . We get two trees,  $\mathcal{T}_a(f)$ , the part which contains  $v_i$  and  $\mathcal{T}_u(f)$  containing  $e_i$ . Then we stick again  $\mathcal{T}_u(f)$  to  $\mathcal{T}_a(f)$ , sticking the edge  $e_i$  on  $v_i$  as an horizontal edge. We add a new vertical edge under  $v_i$  decorated with 1 and ending by an arrowhead decorated with (0). Since  $v_i$  has a valency greater or equal to 3 on  $\mathcal{T}(f)$ , it has a valency greater or equal to 4 on  $\mathcal{T}'(f)$ .

To put f in P-good coordinates we made the choice of a. It is easy to verify that actually the tree doesn't depend on the choice of a. Then in this case we have unicity of  $\mathcal{T}'(f)$  from  $\mathcal{T}(f)$ .

### Sub-Case 2. i = n.

The edge  $e_n$  under the vertex  $v_n$  ends with an arrow decorated with (k). Three cases may arise:

- If  $k \neq 0$  then we are in P-good coordinates.
- If k = 0 and the valency of  $v_n$  is strictly greater than 3 then we are in P-good coordinates.
- If k = 0 and the valency of  $v_n$  is equal 3 then we are not in P-good coordinates.

We study this last case. The face polynomial at  $v_n$  is

$$\boldsymbol{x}^{k'}(z-a\boldsymbol{x}^{\boldsymbol{q}^n})^m.$$

To eliminate the face  $\Gamma_n$ , we have to perform the (unique) change of variables

$$z = z_1' + a \boldsymbol{x}^{\boldsymbol{q}^n}.$$

We considered the Newton map:

$$x = y$$
,  $z = y^{q^n}(z_1 + a)$ 

and we have  $z'_1 = x^{q^n} z_1$ . The computation are the same than in Sub-Case 1.

To get  $\mathcal{T}'(f)$  from  $\mathcal{T}(f)$ , we do the inverse operation than in Sub-Case 1. We delete the edge  $e_n$ . Denote by e the horizontal edge attached to  $v_n$ . We cut  $\mathcal{T}(f)$  in two pieces, separating  $v_n$  and e and we stick it back so that e becomes the vertical edge under  $v_n$ . The vertex  $v_n$  has now valency 2 and has to be eliminated.

Note that we are not necessarily in P-good coordinates yet. We illustrate this fact in several examples:

- In Figure 4(a) we are in suitable coordinates and in Sub-Case 2.
- In Figure 4(b) we have eliminated the face  $\Gamma_n$  but we are not yet in P-good coordinates. We are in Sub-Case 1.
- In Figure 5 we are in P-good coordinates.

We have proven that if f is in suitable coordinates such that its Newton polyhedron is a monotone polygonal path, then there is essentially a unique way to find P-good coordinates.

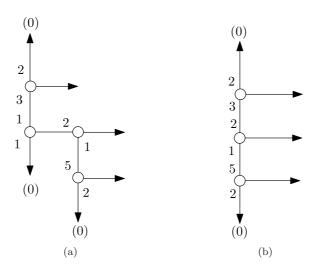


FIGURE 4. Newton tree in suitable coordinates and not P-good

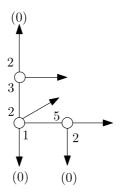


FIGURE 5. Newton tree in P-good coordinates

3.1. On  $\nu$ -quasi-ordinary polynomials whose Newton tree ends only with arrow-heads. Now we assume that f is such that there exists a system of suitable coordinates such that  $\mathcal{T}(f)$  has only arrowheads (no black boxes).

In that case, the Newton polyhedron of f is a monotone polygonal path. We can find a system of P-good coordinates for f. Let  $\mathcal{T}_1(f)$  denote the Newton tree in this system. It has also only arrowheads. Consider the Newton transforms of f in this system of P-good coordinates; their Newton polyhedra are monotone polygonal paths. Then we can find a system of P-good coordinates for them, and so on, see Proposition 3.2. We can summarize it as follows.

**Definition 3.3.** A Newton tree is said to be a P-good tree if the vertices with decoration p = 1 are at the bottom of the tree, connected to a vertical edge ending with an arrowhead decorated with (0), and with valency strictly greater than 3.

**Proposition 3.4.** Starting from a Newton tree  $\mathcal{T}(f)$  (for suitable system of coordinates), we can associate to f a unique P-good tree.

**Question 3.5.** Assume there are two suitable systems of coordinates such that the Newton trees  $\mathcal{T}(f)$  and  $\mathcal{T}'(f)$  have only arrowheads. Do we get the same P-good tree?

In order to answer this question, we distinguish three cases:

- (1) The change of variables doesn't change the Newton diagram of f. It doesn't change either the face polynomials, then  $\mathcal{T}(f) = \mathcal{T}'(f)$ . Then there is one P-good tree.
- (2) The change of variables modifies the Newton diagram as follows: The edges  $\Gamma_1, \ldots, \Gamma_i$  remain unchanged and  $\Gamma_{i+1}, \ldots, \Gamma_n$  are replaced by a new face  $\Gamma$ . In any case  $\mathcal{T}'(f)$  is not in P-good coordinates and we are in Sub-Case 2 in the proof of Proposition 3.2. In order to get P-good coordinates we first have to come back to  $\mathcal{T}(f)$ . Then we have the same P-good tree.
- (3) The change of coordinates changes the Newton diagram as follows: It keeps unchanged  $\Gamma_1, \ldots, \Gamma_i$ , and replace  $\Gamma_{i+1}$  by a face with same equations which hits the hyperplane  $\{z=0\}$ . Either one of  $\mathcal{T}(f)$  or  $\mathcal{T}'(f)$  is P-good and is the P-good tree of the other, or the two trees only differ by exchanging one vertical edge and one horizontal edge. They have the same P-good tree because in the P-good tree they are all horizontal.

Hence, if the Newton tree  $\mathcal{T}(f)$  for f, in suitable coordinates, has no black box then we associate to it a unique P-good tree denoted by  $\mathcal{T}_P(f)$  and the answer to Question 3.5 is positive.

Remark 3.6. A P-good tree is not a minimal tree in the sense of Eisenbud and Neumann [11]. A minimal tree is unique if we forget about the direction of the edges, but not in the strong sense. A P-good tree is unique in the strong sense.

This allows us to define the depth of f in suitable coordinates such that  $\mathcal{T}(f)$  has no black box.

**Definition 3.7.** The *depth* of f is the maximal length of horizontal paths in  $\mathcal{T}_P(f)$ , denoted by depth(f).

## 4. Quasi-Ordinary Polynomials

**Notation 4.1.** Let  $g(x, z) \in \mathbb{K}[[x]][z]$  be a polynomial of degree m and regular of order  $n \leq m$  with coefficients in the formal power series ring  $\mathbb{K}[[x]]$ , that is,

$$g(\mathbf{x}, z) := z^m + a_1(\mathbf{x})z^{m-1} + \ldots + a_{m-n}(\mathbf{x})z^n + a_{m-n+1}(\mathbf{x})z^{n-1} + \ldots + a_{m-1}(\mathbf{x})z + a_m(\mathbf{x}),$$

with  $a_i(\mathbf{0}) = 0$  for  $m - n + 1 \leq i \leq m$  and  $a_{m-n}(\mathbf{0}) \in \mathbb{K}^*$  and let  $f(\boldsymbol{x}, z) := a_0 x_1^{n_1} \cdot \ldots \cdot x_d^{n_d} g(\boldsymbol{x}, z)$ , where  $a_0 \in \mathbb{K}^*$ . Applying Weierstrass Preparation Theorem to  $g(\boldsymbol{x}, z)$  there exists a unit  $u(\boldsymbol{x}, z) \in \mathbb{K}[[\boldsymbol{x}]][z]$  and a Weierstrass polynomial

$$h(x, z) := z^n + B_1(x)z^{n-1} + \ldots + B_{n-1}(x)z + B_n(x)$$

such that  $g(\mathbf{x}, z) = h(\mathbf{x}, z)u(\mathbf{x}, z)$ . In fact  $\deg_z u(\mathbf{x}, z) = m - n$ , see for example [18, Chapter I, p. 11].

**Definition 4.2.** A polynomial  $f \in \mathbb{K}[[x]][z]$  as in Notation 4.1 is called *quasi-ordinary* if its discriminant  $\Delta_z(f) := \operatorname{Res}_z(f, \frac{\partial f}{\partial z})(x)$  is

$$\operatorname{Res}_{z}\left(f, \frac{\partial f}{\partial z}\right)(\boldsymbol{x}) = \boldsymbol{x}^{\boldsymbol{\alpha}}\eta(\boldsymbol{x})$$

with  $\alpha \in \mathbb{N}^d$  and  $\eta(\mathbf{0}) \neq 0$ .

Remark 4.3. Since  $\operatorname{Res}_z(f, f_1 f_2) = \operatorname{Res}_z(f, f_1) \operatorname{Res}_z(f, f_2)$ , a polynomial  $f \in \mathbb{K}[[x]][z]$  as in Notation 4.1 is quasi-ordinary if and only if g is quasi-ordinary.

If  $f = f_1 f_2 \in \mathbb{K}[[x]][z]$  then  $f_i$  is also quasi-ordinary. It follows from the factoriality of the ring  $\mathbb{K}[[x]]$  and the well-known property of discriminants  $\Delta_z(f_1 f_2) = \Delta_z(f_1)\Delta_z(f_2)(\operatorname{Res}_z(f_1, f_2))^2$ .

In this section we prove the following theorem.

**Theorem 4.4.** Let  $f \in \mathbb{K}[[x]][z]$  be a polynomial as in Notation 4.1. Then f is quasi-ordinary if and only if there exists a suitable system of coordinates of f such that its Newton tree  $\mathcal{T}(f)$  has only arrow-heads decorated with (0) and (1) and has no black boxes.

Quasi-ordinary power series were introduced by Zariski using the discriminant. In fact quasi-ordinary power series are the natural generalization, in the sense of the factorization theorem given by Jung and Abhyankar [1], of algebraic curves.

The fact that at each stage of the Newton process, using eventually automorphisms, its Newton polyhedron is a monotone polygonal path is very useful. It is one of the main ingredients in the proof in [4] of the monodromy conjecture for hypersurfaces defined by quasi-ordinary power series in arbitrary dimension.

First we prove that if  $f \in \mathbb{K}[[x]][z]$  is quasi-ordinary, the Newton process ends with monomials multiplied by a unit.

**Proposition 4.5** [21]. Let  $f \in \mathbb{K}[[x]][z]$  be a quasi-ordinary polynomial, then there exists a power series  $b(x) \in \mathbb{K}[[x]]$  such that  $f(x, z - b(x)) \in \mathbb{K}[[x]][z]$  is a  $\nu$ -quasi-ordinary polynomial.

**Lemma 4.6** [16, Lemma 3.16]. If f is a quasi-ordinary polynomial in  $\mathbb{K}[[x]][z]$ , then there exists a system of coordinates such that its Newton polyhedron  $\mathcal{N}(f)$  is a monotone polygonal path.

*Proof.* Since  $f(\mathbf{x}, z) = \mathbf{x}^n h(\mathbf{x}, z) u(\mathbf{x}, z)$  the Newton diagram  $\mathcal{N}_+(f)$  is the Minkowski sum of each of the Newton diagram of its factors it is enough to prove for h which is regular of order n in z.

We will work by induction on the order of h in z. It is true for n=1. We assume that it is true for all regular polynomial of order strictly less than n. Since h is a regular quasi-ordinary polynomial, by Proposition 4.5, there exists a change of coordinates such that h is  $\nu$ -quasi-ordinary with a  $\nu$ -proper face  $\Gamma_1$ . By Factorization Theorem 1.13 and Corollary 1.16 there exists an elementary polynomial  $f_1$  in  $\mathbb{K}[[x]][z]$  with Newton diagram parallel to  $\Gamma_1$  and a polynomial q(x, z) in  $\mathbb{K}[[x]][z]$ 

such that  $h = f_1q$ . Since h is quasi-ordinary and q is one of its factors then q is a quasi-ordinary polynomial of order strictly less than n. From the hypothesis, there exists a system of coordinates such that  $\mathcal{N}(q)$  is a monotone polygonal path, and this change of coordinates does not change the Newton diagram of  $f_1$ . Using the fact that the Newton diagram of a product is the Minkowski sum of the Newton diagrams of each factors, we deduce that the Newton polyhedron of h is a monotone polygonal path.

In particular there exist P-good coordinates for quasi-ordinary polynomials f.

**Lemma 4.7.** [4, Chapter 3, Lemma 3.21] If f is quasi-ordinary, after a Newton map  $\sigma_{\Gamma,j}$ , the total transform  $f_{\Gamma_1,j} := \sigma_{\Gamma_1,j}(f) \in \mathbb{K}[[\boldsymbol{y}]][z_2]$  is a quasi-ordinary polynomial in  $\mathbb{K}[[\boldsymbol{y}]][z_2]$ .

Furthermore the depth of a quasi-ordinary polynomial f in P-good coordinates was defined similarly as in [4, Chapter 3, Definition 24], and it decreases after Newton maps.

**Lemma 4.8.** Let  $f \in \mathbb{K}[[x]][z]$  be a polynomial as in Notation 4.1. If f is quasi-ordinary then there exists a system of suitable coordinates such that its Newton tree  $\mathcal{T}(f)$  has only arrow-heads decorated with (0) and (1) and has no black boxes.

*Proof.* In fact in [4, Chapter 3, Section 3.26] it was defined Newton tree  $\mathcal{T}(f)$  for of a quasi-ordinary polynomials f in P-good coordinates and in this tree all ends are arrow-heads and if an arrow-head is not a dead-end then it is decorated with (1).  $\square$ 

To prove the converse it is enough to prove that  $f \in \mathbb{K}[[x]][z]$  and its polar  $f_z := \frac{\partial f}{\partial z} \in \mathbb{K}[[x]][z]$  are separated polynomials because after Theorem 2.20 they are comparable series.

We assume that there is a system of suitable coordinates such that  $\mathcal{T}(f)$  has no black box and arrowheads decorated with (0) and (1). We consider coordinates such that the Newton tree is  $\mathcal{T}_P(f)$ . Since P-good coordinates are suitable coordinates, these coordinates are suitable for  $ff_z$ .

that the Newton Good Z, these coordinates are suitable for  $ff_z$ . Given a linear form D given by  $\sum_{j=1}^d a_j\alpha_j + b\beta$ ,  $a_j$ ,  $b \in \mathbb{N}$ , we define  $f_D(\boldsymbol{x}, z)$  as the sum of the monomials  $c_{\boldsymbol{\alpha},\beta}\boldsymbol{x}^{\boldsymbol{\alpha}}z^{\beta}$ ,  $c_{\boldsymbol{\alpha},\beta} \neq 0$ , for which  $D(\boldsymbol{\alpha},\beta)$  is minimal. Note that  $f_D$  is not a monomial if and only there is a face of the Newton polyhedron contained in an affine hyperplane defined by D.

**Lemma 4.9.** Assume that D is a linear form such that

$$f_D(\boldsymbol{x},\,z) = \boldsymbol{x}^{\boldsymbol{A}} z^B,$$

with  $B \neq 0$ , then there is no face of  $\mathcal{N}(f_z)$  contained in an affine hyperplane defined by D.

*Proof.* Denote by  $D_+$  the space limited by the first quadrant and the affine hyperplane defined by D which hits the Newton polyhedron of f and doesn't contain the

origin. We have

$$f(\boldsymbol{x}, z) = f_D(\boldsymbol{x}, z) + \sum_{(\boldsymbol{\alpha}, \beta) \in D_+, \, \beta > 0} c_{\boldsymbol{\alpha}, \beta} \boldsymbol{x}^{\boldsymbol{\alpha}} z^{\beta} + \sum_{(\boldsymbol{\alpha}, \beta) \in D_+, \, \beta = 0} c_{\boldsymbol{\alpha}, \beta} \boldsymbol{x}^{\boldsymbol{\alpha}} z^{\beta},$$
$$f_z(\boldsymbol{x}, z) = B \boldsymbol{x}^{\boldsymbol{A}} z^{B-1} + \sum_{(\boldsymbol{\alpha}, \beta) \in D_+, \, \beta > 0} \beta c_{\boldsymbol{\alpha}, \beta} \boldsymbol{x}^{\boldsymbol{\alpha}} z^{\beta-1}.$$

Then, there is no face of the Newton diagram of  $f_z$  in an affine hyperplane defined by D.

**Lemma 4.10.** Assume  $\Gamma_{\ell}$  is a face of the Newton diagram of f which doesn't hit the hyperplane  $\{z=0\}$  with face polynomial, see (1.6),

$$f_{\Gamma_\ell} = a_{\Gamma_\ell} oldsymbol{x}^{oldsymbol{n}^\ell} z^{n_{d+1}^\ell} \prod_{i=1}^{k^\ell} (z^{p_\ell} - \mu_j^\ell oldsymbol{x}^{oldsymbol{q}^\ell})^{m_j^\ell},$$

the factors being irreducible in  $\mathbb{K}[[x]][z]$ , i.e.,  $\gcd(q^{\ell}, p_{\ell}) = 1$ ,  $\mu_j^{\ell} \in \mathbb{K}^*$  with  $\mu_j^{\ell} \neq \mu_i^{\ell}$  and  $a_{\Gamma_{\ell}} \in \mathbb{K}^*$ .

Then the Newton diagram of  $f_z$  has a face  $\Gamma'_{\ell}$  parallel to  $\Gamma_{\ell}$  and the corresponding initial form  $(f_z)_{\Gamma'_{\ell}}$  is

$$a_{\Gamma_{\ell}} \boldsymbol{x}^{\boldsymbol{n}^{\ell}} z^{n_{d+1}^{\ell} - 1} \left( \prod_{j=1}^{k^{\ell}} (z^{p_{\ell}} - \mu_{j}^{\ell} \boldsymbol{x}^{\boldsymbol{q}^{\ell}})^{m_{j}^{\ell} - 1} \right) \times \left( \sum_{i=1}^{k^{\ell}} \left( (n_{d+1}^{\ell} + p_{\ell} m_{i}^{\ell}) z^{p_{\ell}} - n_{d+1}^{\ell} \mu_{i}^{\ell} \boldsymbol{x}^{\boldsymbol{q}^{\ell}} \right) \prod_{j \neq i} (z^{p_{\ell}} - \mu_{j}^{\ell} \boldsymbol{x}^{\boldsymbol{q}^{\ell}}) \right).$$

*Proof.* Same computation as before.

**Lemma 4.11.** Assume  $\Gamma_{\ell}$  is the face of the Newton diagram of f which hits the hyperplane  $\{z=0\}$  with face polynomial

$$f_{\Gamma_{\ell}} = a_{\Gamma_{\ell}} \boldsymbol{x}^{\boldsymbol{n}_{\ell}} \prod_{j=1}^{k^{\ell}} (z^{p_{\ell}} - \mu_{j}^{\ell} \boldsymbol{x}^{\boldsymbol{q}_{\ell}})^{m_{j}^{\ell}}.$$

Then if  $k^l > 1$  or if  $k^l = 1$  and  $m_1^l > 1$  the Newton diagram of  $f_z$  has a face parallel to  $\Gamma_\ell$ , denoted by  $\Gamma'_\ell$  and

$$(f_z)_{\Gamma'_{\ell}} = p_{\ell} a_{\Gamma_{\ell}} \boldsymbol{x}^{n_{\ell}} z^{p_{\ell} - 1} \prod_{j=1}^{k^{\ell}} (z^{p_{\ell}} - \mu_j^{\ell} \boldsymbol{x}^{q_{\ell}})^{m_j^{\ell} - 1} \left( \sum_{i=1}^{k^{\ell}} m_i^{\ell} \prod_{j \neq i} (z^{p_{\ell}} - \mu_j^{\ell} \boldsymbol{x}_{\ell}^{q}) \right).$$
(4.1)

If  $k^{\ell} = 1$  and  $m_1^{\ell} = 1$  then the Newton diagram of  $f_z$  has no face parallel to  $\Gamma_{\ell}$ .

Remark 4.12. If  $k^{\ell} = 1$  and  $m_1^{\ell} = 1$  then  $p_{\ell} > 1$  since we are in P-good coordinates for f. Then either  $f_z$  is divisible by  $z^{p_{\ell}-1}$  or  $\mathcal{N}(f_z)$  has faces which are not faces of f. The proof of the above lemma is straightforward.

**Corollary 4.13.** Consider the bicolored Newton tree  $\mathcal{T}(ff_z)$  colored blue for f and red for  $f_z$ . Denote by  $v_1, \ldots, v_n$  the vertices on the first vertical line.

There exists  $n_0$ ,  $1 \le n_0 \le n$ , such that for  $1 \le j \le n_0$  there is a blue or blue-red horizontal edge attached to  $v_j$  and for  $n_0 < j \le n$  there is no blue, neither blue-red horizontal edge attached to  $v_j$ .

Attached to  $v_j$ ,  $1 \leq j \leq n_0$ , there are eventually blue horizontal edges ending with an arrow and/or blue-red horizontal edges ending with a vertex and always red horizontal edges.

The horizontal edges corresponding to the roots coming from the factor

$$\sum_{i=1}^{k^{\ell}} \left( (n_{d+1}^{\ell} + p_{\ell} m_i^{\ell}) z^{p_{\ell}} - n_{d+1}^{\ell} \mu_i^{\ell} \boldsymbol{x}^{\boldsymbol{q}^{\ell}} \right) \prod_{j \neq i} (z^{p_{\ell}} - \mu_j^{\ell} \boldsymbol{x}^{\boldsymbol{q}^{\ell}}),$$

of  $(f_z)_{\Gamma'_\ell}$  are red-colored. The degree in z is  $k^\ell p_\ell$ .

Proof of the implication  $\Leftarrow$  in Theorem 4.4. We use induction on  $\delta := \text{depth}(f)$ . Let us start with the case  $\delta = 0$ . We have  $f = x^n z$  up to a unit and  $f_z = x^n$ . Then f and  $f_z$  are separated.

For  $\delta = 1$ , consider a vertex  $v_j$  of the tree of  $ff_z$ . If  $1 < j \le n_0$ , there are exactly  $k_j$  blue horizontal edges attached to  $v_j$  ending by an arrow, and some red edges, but no blue-red edges. For  $n_0 < j \le n$ , there are only red edges. Then f and  $f_z$  are separated.

Assume that for any g in P-good coordinates with  $\mathcal{T}_P(g)$  with no black box and arrowheads decorated with (0) and (1) of depth  $\delta' < \delta$ , g and  $g_z$  are separated.

Consider f in P-good coordinates of depth  $\delta$  and assume that  $\mathcal{T}(f)$  has no black box and has arrowheads decorated with (0) and (1). If  $\delta > 1$  there is a vertex  $v_j$ ,  $1 \leq j \leq n_0$  such that there is a blue-red horizontal edge attached to  $v_j$ . We consider the corresponding Newton map  $\sigma$ . Consider the polynomials  $(f_z)_{\sigma}(\boldsymbol{y}, z_1)$  and  $f_{\sigma}(\boldsymbol{y}, z_1)$ , after the change of variables to have  $f_{\sigma}(\boldsymbol{y}, z_1)$  in P-good coordinates. Using see Lemma 1.18 we have  $(f_{\sigma})_{z_1} = \boldsymbol{x}^q(f_z)_{\sigma}$ .

Therefore,  $\mathcal{T}(f_{\sigma} \cdot (f_z)_{\sigma})$  and  $\mathcal{T}(f_{\sigma} \cdot (f_{\sigma})_{z_1})$  are the same except for the decoration of the top arrow. Since  $\operatorname{depth}(f_{\sigma}) < \delta$ , then  $f_{\sigma}$  and  $(f_{\sigma})_{z_1}$  are separated on  $\mathcal{T}(f_{\sigma} \cdot (f_{\sigma})_{z_1})$ . The tree  $\mathcal{T}(f \cdot f_z)$  is obtained by gluing  $\mathcal{T}(f_{\sigma} \cdot (f_z)_{\sigma})$  on  $v_j$ . We can apply that to every vertex where we have a blue-red edge attached. Then f and  $f_z$  are separated.

This description allows us to study how f and its derivative  $f_z$  separate on the Newton tree of f.

**Definition 4.14.** The edges at the bottom of the polyhedron ending with an arrowhead of multiplicity (0) (dead ends) are called *leaves*, and the vertices at the other end of the arrow-head *leaf vertices*.

**Theorem 4.15.** Let  $f \in \mathbb{K}[[x]][z]$  be a polynomial as in Notation 4.1 in a P-good system of coordinates and we consider the Newton tree  $\mathcal{T}(f)$  whose first vertical line has  $v_1, \ldots, v_s$  as vertices.

(D1) The polar  $f_z$  doesn't separate on the edges of the Newton tree of f which are not leaves of f.

- (D2) On each vertex  $v_{\ell}$  of the Newton tree of f, different from a leaf vertex, the polar  $f_z$  separates with order equal to  $k^{\ell}p_{\ell}$ , see (4.1); recall  $k^{\ell}+2$  is the valency of the vertex and  $p_{\ell}$  the decoration under the vertex.
- (D3) The polar  $f_z$  separates on a leaf vertex or on a leaf with total order  $\frac{k^s-1}{p_s} + p_s 1$

*Proof.* It follows from the previous discussion.

Remark 4.16. We don't know anything on  $f_z$  after its separation with f. In particular in general the derivative is not quasi-ordinary. In the case d = 1 this gives Lê-Michel-Weber Theorem [22].

We can compute the discriminant of f from its Newton tree. The following result follows from a recursive application of [5, Lemmas 5.7, 5.10], Proposition 2.21 and Theorem 4.15.

**Proposition 4.17.** Let f be a quasi-ordinary Weierstrass polynomial of degree n. Its discriminant can be computed from the polyhedron of f by

$$\Delta_z(f) := \operatorname{Res}_z(f, f_z)(\boldsymbol{x}) = \boldsymbol{x}^D u(\boldsymbol{x}), \quad \text{with } u(0) \neq 0,$$

where

$$oldsymbol{D} := (D_1, \, \dots, \, D_d) = \sum_v (\delta_v - 2) oldsymbol{N}_v' - \sum_{v \text{ leaf vertex}} rac{oldsymbol{N}_v'}{p_v},$$

where  $\delta_v$  is the valency of the vertex v, and  $\mathbf{N}'_v = \boldsymbol{\rho} * \mathbf{N}_v$ , where  $\boldsymbol{\rho} := \prod_{w \text{ before } v} \boldsymbol{c}_w$  (recall that  $\boldsymbol{c}_w$  is the gcd's of the vertex w).

This can be seen as a generalisation of Kouchnirenko theorem [20], [9] and as a generalization of [13, Corollaire 5], where P. González Pérez gave information about the Newton polyhedron of the resultant of two quasi-ordinary hypersurfaces satisfying an appropriate non-degeneracy condition.

#### 5. Transversal sections

Consider a quasi-ordinary polynomial  $f \in \mathbb{K}[[x]][z]$ . Fix  $i \in \{1, \ldots, d\}$ . We follow this convention: if p is a d-tuple, then  $p^i$  is the (d-1)-tuple  $(p_1, \ldots, p_{i-1}, p_{i+1}, \ldots, p_d)$ . Let  $K_i$  be an algebraic closure of  $K((x_i))$ . The i-transversal section of h is the quasi-ordinary polynomial  $h^i \in K_i[[x^i]][z]$  obtained from h considering  $x_i$  as a generic constant. For  $I \subset \{1, \ldots, d\}$  we denote  $I' := \{1, \ldots, d\} \setminus I$ ; we can define recursively the I-transversal section  $h^I$  and if I' has one element, say j, we call it j-curve transversal section  $h^{I_j}$ .

Consider f a quasi-ordinary polynomial in suitable coordinates such that its Newton tree has no black box. We first study the Newton trees of the transversal sections in the same system of coordinates.

Then we assume the Newton tree of f has only one arrow-head with positive multiplicity. We show that we can retrieve its decorated Newton tree from the Newton tree of its curve transversal sections in the same system of coordinates.

We give examples where this is no more true when the Newton tree of f has more than one arrow-head with positive multiplicity. Nevertheless we can prove that we

can retrieve the global decorations of the vertices of f from the decorations of the curve transversal sections. This result was a crucial ingredient in the proof of the monodromy conjecture for quasi-ordinary singularities in arbitrary dimension [4].

Let f be a quasi-ordinary polynomial in suitable coordinates such that the Newton tree  $\mathcal{T}(f)$  has no black box. We are going to describe how we can find the tree of the *i*-transversal section  $f^i$  in the same coordinates.

**5.1. Construction of**  $\mathcal{T}_{\mathcal{N}}(f^i)$ . The general procedure is that we can copy  $\mathcal{T}_{\mathcal{N}}(f)$ and erase the i-th component of the local numerical data except in 3 cases which we develop below.

Case 1. Along  $\mathcal{T}_{\mathcal{N}}(f)$ , there is a sequence of consecutive vertices where the local numerical data coincide after erasing the *i*-th coordinate (and dividing by the gcd). This means that the two faces corresponding to these vertices project on the same face. Then the chain of vertices project on  $\mathcal{T}_{\mathcal{N}}(f^i)$  on the same vertex decorated with the common normalized numerical data.

Case 2. The local numerical data at v on  $\mathcal{T}_{\mathcal{N}}(f)$  satisfy  $\gcd(q^i, p) =: d^i$ , with  $d^i > 1$ , and let  $q'_j := \frac{q_j}{d^i}$ ,  $j \in I'$ , and  $p' := \frac{p}{d^i}$ . Assume that the face polynomial at v is

$$\prod_{j} (z^p - \mu_j \boldsymbol{x}^{\boldsymbol{q}})^{m_j}.$$

On  $\mathcal{T}_{\mathcal{N}}(f^i)$ , there is a vertex  $v^i$  decorated with (q', p') and with face polynomial

$$\prod_{j} \prod_{\{\zeta : \zeta^{d^i} = 1\}} (z^{p'} - \zeta \mu_j x_i^{\frac{q_i}{d^i}} (\boldsymbol{x}^i)^{\boldsymbol{q}'})^{m_j}.$$

Then if k horizontal edges arise from v,  $kd^i$  arise from  $v^i$ .

Remark 5.1. Case 1 and 2 are not exclusive. The number of horizontal edges at  $v^i$  is obtained as the sum, along the vertices v projecting on  $v^i$ , of  $d^i(v)$  times the number of horizontal edges arising from v.

**Example 5.2.** We have in Figure 6(a) the tree of

$$f(x_1, x_2, z) := (z^2 - x_1^3 x_2)(z^2 - x_1^3 x_2^4)(z^2 - x_1^5 x_2^6) \in K[[x_1, x_2]][z];$$

Figures 6(b) and 6(c) show the trees for  $f(x_1, x_2, z) \in K_1[[x_2]][z]$  and  $f(x_1, x_2, z) \in K_1[[x_2]][z]$  $K_2[[x_1]][z]$ , respectively. This example illustrates Cases 1 and 2. More explicitely, the tree for  $K_2[[x_1]][z]$  illustrates Case 1 and the tree for  $K_1[[x_2]][z]$  illustrates

Case 3. We have  $q_{\ell}^i = \mathbf{0}^i$  for the first h vertices of  $\mathcal{T}_{\mathcal{N}}(f)$ . That means that the projection of the face  $\Gamma_{\ell}$  of  $\mathcal{N}(f)$  has equation  $px_j = 0, j \in I'$ ; the polynomial  $f_{\Gamma_{\ell}}(\boldsymbol{x}, z)$  doesn't depend on  $x_j, j \in I'$ , that is,

$$f_{\Gamma_{\ell}}(\boldsymbol{x}^{i},\,z) = z^{k_{\ell}}\boldsymbol{x}^{N} \prod_{j=1}^{r_{\ell}} (z^{p_{\ell}} - \mu_{j}^{\ell}x_{i}^{q_{\ell}})^{m_{j}^{\ell}},$$

i.e.,  $f^i$  decomposes in  $K_i[[x^i]][z]$  in  $m := \sum_{\ell=1}^h p_\ell \sum_{j=1}^{r_\ell} m_j^\ell$  factors based at the  $p_{\ell}$ -roots  $\mu_j^{\ell} x_i^{q_{\ell}}$  of z-degree  $m_j^{\ell}$ , and one extra factor based at z=0 of degree  $k_h$  (if

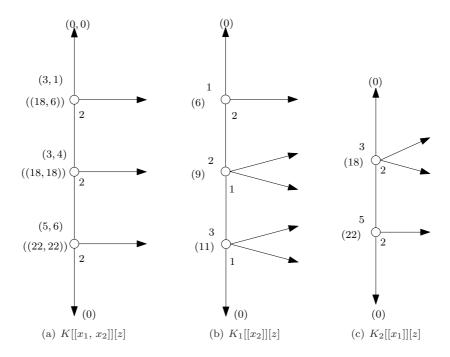


FIGURE 6. Newton trees of Example 5.2

 $k_h > 0$ ). We proceed as follows. We keep m void trees and if  $k_h > 0$  we consider also a tree for z = 0, where we erase the vertices  $v_1, \ldots, v_h$  and we keep the upper arrow.

**5.2.** Construction of  $\mathcal{T}(f^i)$  without decorations. Once we have obtained the vertical Newton trees we explain how to obtain the trees  $\mathcal{T}^i(f)$ , when we see  $f \in K_i[[\boldsymbol{x}^i]][z]$ . We proceed by induction of depth(f). If the depth is one we have  $\mathcal{T}^i(f) = \mathcal{T}^i_{\mathcal{N}}(f)$  with the following exception. If we are in Case 3, we replace the empty trees by the ones in Figure 1(a).

If the depth is greater than one, we start with  $\mathcal{T}_{\mathcal{N}}(f^i)$ . Each arrow of this tree is related to one Newton transformation  $\sigma := \sigma_{\Gamma,\mu}$  and by the induction hypothesis we may assume that  $\mathcal{T}(f^i_{\sigma})$  are constructed  $\forall \Gamma, \mu$ . In order to recover  $\mathcal{T}(f^i)$  we proceed as in Step 3 of Section 2.1 with the following caveats.

If we are not in Case 3 and  $\mathcal{T}(f^i_\sigma)$  is a disjoint union of  $k_{\Gamma,\mu}$  trees, then we perform Step 3 for each connected component. Hence we will obtain  $\prod_{\Gamma,\mu} k_{\Gamma,\mu}$  disconnected trees.

If we are in Case 3, we proceed in the same way for the tree corresponding to z=0 (if  $k_h>0$ ). Each void tree is associated to a pair  $(\Gamma_\ell, \mu_j^\ell)$ ,  $\ell \leq h$ , and we simply take  $p_\ell$  copies of the tree  $\mathcal{T}(f_\sigma^i)$ ,  $\sigma = \sigma_{\Gamma_\ell, \mu_j^\ell}$ .

Remark 5.3. In order to construct Newton maps, we followed some conventions in Section 1.3 which may not pass to transversal sections. It is easily seen that the results do not depend on the particular choice of conventions. In particular,

everything works if one does not choose a tuple u in Section 1.3 but a p-th root of  $\mu_i$  instead.

**5.3.** Decorations of  $\mathcal{T}(f^i)$ . In order to finish the construction of the Newton trees  $\mathcal{T}^i(f)$  we need to compute the decorations. In order to do it in a simpler way we define new decorations to the Newton tree.

**Definition 5.4.** The prime local decorations  $\langle R_1^v, \ldots, R_d^v; p^v \rangle$  of a vertex v are defined as follows. They coincide with the local decoration if v is in the first vertical tree. If not, they are computed with the following recursive formula (as in Lemma 2.1), where w is the preceding vertex:

$$R_{j}^{v} := q_{j}^{v} + \frac{p^{v} R_{j}^{w} p^{w}}{\gcd(R_{j}^{w}, p^{w})^{2}}.$$

Note that it is possible to compute local decorations from prime local decorations (and viceversa). Moreover, they coincide if the decorations are *coprime enough*, e.g., in the curve case. The reason to define these decorations is that they behave better when passing to transversal section since they really depend on the quotients  $\frac{q_j}{p}$  and not on the pairs  $(q_j, p)$ .

In order to decorate the Newton trees of the transversal sections the strategy is as follows:

(TSNT1) In  $\mathcal{T}(f)$  compute the prime local decorations from the local decorations.

(TSNT2) Pick a vertex  $v^i$  in  $\mathcal{T}^i(f)$ , and consider the prime local decorations of the vertex v in  $\mathcal{T}(f)$  which originates  $v^i$ . Forget the i-th coordinate and make the decoration coprime as in Case 2.

(TSNT3) Obtain the local decoration of  $v^i$  from the prime local decorations.

**5.4. Examples of transversal sections.** We illustrate the above theory with some examples.

**Example 5.5.** The tree of Figure 7(a) corresponds to

$$f(x_1, x_2, x_3, z) = (z^7 - x_1^2)^2 (z^3 - x_1^5 x_2 x_3) + x_1^{10} x_2 x_3 \in K[[x_1, x_2, x_3]][z]$$

and illustrates Case 3. If we consider  $f \in K_1[[x_2, x_3]][z]$ , we have  $f(x_1, 0, 0, z) = (z^7 - x_1^2)^2 z^3$ , i.e,  $f = f_0 \prod_{j=1}^7 f_j$ , where  $f_0(x_1, 0, 0, 0) = 0$  and  $\{f_j(x_1, 0, 0, 0)\}_{j=1}^7 = \{b \in K_1: b^7 = x_1^2\}$ . Hence we obtain 8 disjoint Newton trees, one for  $f_0$  and 7 equal trees for  $f_j$ ,  $j = 1, \ldots, 7$ .

These following examples illustrate Case 3 at the second step of the algorithm.

**Example 5.6.** Figure 8(a) shows the tree of

$$f(x_1, x_2, z) = ((z^7 - x_1^2 x_2^3)^2 - x_1^5 x_2^6)^2 + x_1^{11} x_2^{13} \in K[[x_1, x_2]][z],$$

while Figure 8(b) shows the one for  $f \in K_1[[x_2]][z]$ .

Example 5.7. Figure 9(a) shows the tree of

$$f(x_1, x_2, z) = (z^2 - x_1^2 x_2^3)^6 + (z^2 - x_1^2 x_2^3)^3 x_1^7 x_2^9 + x_1^{15} x_2^{19} \in K[[x_1, x_2]][z].$$

while Figure 9(b) shows the one for  $f \in K_1[[x_2]][z]$ .

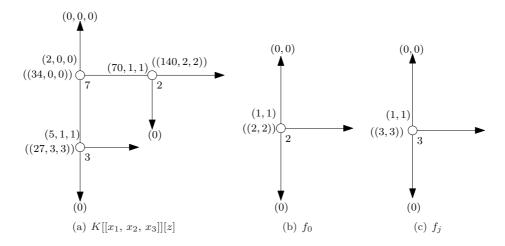


FIGURE 7. Newton trees of Example 5.5.

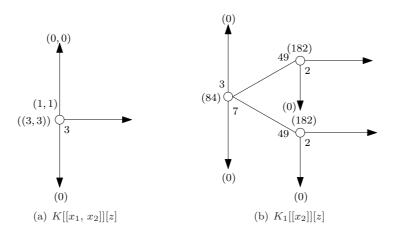


FIGURE 8. Newton trees of Example 5.6.

**5.5.** Curve transversal sections. We summarize the previous process showing the computations for curve transversal sections. Now we are describing the Newton trees of the curve transversal sections in the case where, in suitable coordinates, the Newton tree of f has only one arrow-head with positive multiplicity and no black box. Consider the first vertex and fix  $i \in \{1, \ldots, d\}$ . There are three cases.

First Vertex 1.  $q_i \neq 0$ ,  $gcd(q_i, p) = 1$ .

In this case, we have a first vertex decorated with  $(q_i, p)$  and with a unique horizontal edge starting from this first vertex.

First Vertex 2.  $q_i \neq 0$ ,  $gcd(q_i, p) =: c_i > 1$ .

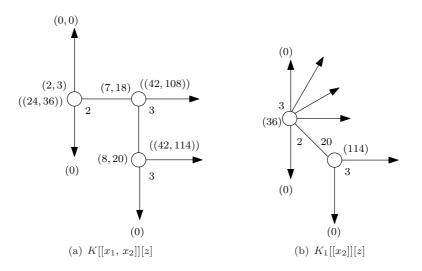


FIGURE 9. Newton trees of Example 5.7.

We have a first vertex decorated with  $(q'_i, p_i)$ ,  $q'_i := \frac{q_i}{c_i}$ ,  $p_i := \frac{p_i}{c_i}$ , and  $c_i$  horizontal edges starting from the vertex.

## First Vertex 3. $q_i = 0$ .

We are not studying the transversal section at the origin for z, but at some other points. There are p of them. In this case we have p Newton trees which begin eventually with the next vertex (as far as this situation is not repeated again).

Now we apply a Newton map  $\sigma_{\Gamma,\mu}$  to go to the following vertex whose decorations in  $\mathcal{T}_{\mathcal{N}}(f_{\Gamma,\mu})$  are  $((r_1,\ldots,r_n),p^1)$ . The local numerical data  $(Q_1,\ldots,Q_d,p^1)$  and the prime local numerical data  $\langle R_1,\ldots,R_d;p^1\rangle$  of the second vertex satisfy

$$Q_i = p^1 p \frac{q_i}{c_i} + r_i, \quad R_i = p^1 p \frac{q_i}{c_i^2} + r_i.$$

Again we have three cases.

# Second Vertex 1. $gcd(Q_i, p^1) = 1, \forall i.$

For all the transversal sections with a first vertex decorated with  $(q_i^t, p^t)$  we add a new vertex to each of the edges decorated with  $(r_i + q_i^t p^t p^1, p^1)$  and we add a unique edge to this vertex. If we do not have a first vertex (the case of First Vertex 3) we begin the tree with a first vertex decorated with  $(r_i, p^1)$  and one horizontal edge starting from the vertex.

**Second Vertex 2.** For some 
$$i$$
,  $gcd(Q_i, p^1) =: c_i^1 > 1$  and  $r_i \neq 0$ .

For the corresponding transversal section, if we have a first vertex decorated with  $(q_i^t, p^t)$  we add a new vertex decorated  $(r_i + p^t p^1 \frac{q_i^t}{c_i}, \frac{p^1}{c_i^1})$  to all the edges and  $c_i^1$  horizontal edges starting from the vertex. If we do not have a first vertex we

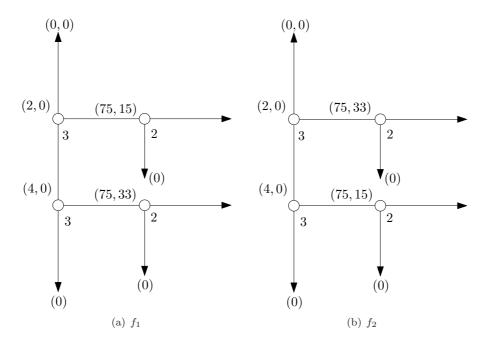


FIGURE 10.

begin the tree with a first vertex decorated with  $(\frac{r_i}{c_i^1}, \frac{p^1}{c_i^1})$  and we add  $c_i^1$  horizontal edges.

**Second Vertex 3.** For some i we have  $r_i = 0$ .

If we have a first vertex decorated with  $(q_i^t, p^t)$ , we stay on the same vertex. Then if we had already p edges we should have  $pp^1$  edges starting from the vertex. If we don't have already any vertex then we will have  $pp^1$  trees starting eventually with the next vertex.

The global process continue with all their corresponding Newton maps. Note that in the curve case prime local decorations and local decorations coincide.

**5.6.** Obtaining  $\mathcal{T}(f)$  from  $\mathcal{T}(f^{I_i})$ . Now we will see how one can retrieve the quasi-ordinary singularity from the transversal sections. This is not always possible as shown in the following examples.

Example 5.8. The quasi-ordinary polynomials

$$f_1 = ((z^3 - x_1^2)^2 + x_1^{25}x_2^{11})((z^3 - x_1^4)^2 + x_1^{25}x_2^5),$$
  

$$f_2 = ((z^3 - x_1^2)^2 + x_1^{25}x_2^5)((z^3 - x_1^4)^2 + x_1^{25}x_2^{11})$$

have the same transversal sections but they have different decorated Newton trees (Figure 10).

**Example 5.9.** Let  $f_n = z^n - x_1 x_2$ . If  $n_1 \neq n_2$  then the decorated Newton tree of  $f_{n_1}$  is not the same than the decorated Newton tree of  $f_{n_2}$ . If we do not keep

the system of coordinates, i.e., exchanging the roles of z and  $x_i$  in each transversal section, we get the same trees, see Figure 11, for any n.

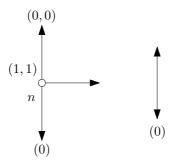


Figure 11.

Example 5.10. The transversal sections of the quasi-ordinary polynomial

$$f_1 = ((z^2 - x_1^3 x_2)^2 + x_1^5 x_2^3 z)((z^2 - x_1^3 x_2^4)^2 + x_1^6 x_2^9 z),$$
  

$$f_2 = ((z^2 - x_1^3 x_2^4)^2 + x_1^5 x_2^9 z)((z^2 - x_1^3 x_2)^2 + x_1^6 x_2^3 z)$$

have the same decorated Newton trees but the two germs have different decorated Newton trees, see Figure 12.

**Definition 5.11.** The *multiplicity* of a tree is the sum over all vertices v of the number of horizontal edges arising from v.

**Theorem 5.12.** Assume that f is a quasi-ordinary polynomial in  $\mathbb{K}[[x]][z]$  whose Newton tree in suitable coordinates has only one arrow-head with positive multiplicity and no black box. Then, there is a unique way to recover the decorated Newton tree of f from the decorated Newton trees, in the same system of coordinates, of all its curve transversal sections.

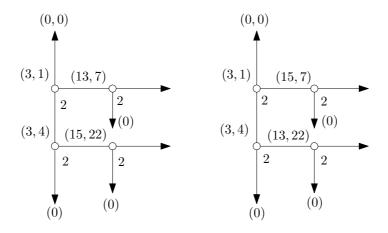


Figure 12.

*Proof.* The proof can be done by induction on the sum of the multiplicities of the curve transversal sections. If this sum is equal to one, it means that we may assume  $f_{\Gamma}(\mathbf{x}, z) = \mathbf{x}^{N}(z^{p} - x_{1}^{q})$  and the result follows immediately.

Let us assume that the result is true when the sum of the multiplicities is less than n, and assume this sum for f equals n > 1. In this case

$$f_{\Gamma}(\boldsymbol{x}, z) = \boldsymbol{x}^{N}(z^{p} - \boldsymbol{x}^{q})^{m}$$

for  $q_1 \ge \dots q_k > q_{k+1} = \dots = q_d = 0$ ; recall that f has only one arrow-head with positive multiplicity and no black box. Then we must have either  $m \ge 2$  or  $k \ge 2$ .

If k = 1 the results follows easily by induction. Let us consider the case  $k \ge 2$ . As we have seen before, after the Newton map we get in general bunches of disconnected trees. Moreover, all the trees in the bunch corresponding to the *i*-th coordinate are equal, so we need only to retain one of them and the number of trees in the bunch.

We start the construction of the tree of f from the trees of the curve transversal sections. We consider all curve transversal sections with a complete tree. There is at least one. We consider the first vertices of these trees. We consider the decorations  $(q_i, p_i)$  of these vertices. Let  $p = \max\{p_1, \ldots, p_d\}$ . We know that  $p_i$  divides p. Let  $p = p_i c_i$ . Then the first vertex of the tree of f is decorated by  $((q_1 c_1, \ldots, q_k c_k, 0, \ldots), p)$ .

We consider the Newton tree of the transforms of the curve transversal sections of f. The sum of the multiplicities is strictly less than n and hence we can apply induction hypothesis and all the results in this sections.

Then we can reconstruct our decorated tree from the transversal sections. There are many branches in the transversal sections, but because they are all the same there is a unique way to reconstruct the tree of f.

Now we prove that we have all the information of the global numerical data can be found in the trees of the transversal sections.

**Proposition 5.13.** For each vertex v of the tree of a quasi ordinary polynomial, and each i = 1, ..., d, the terms  $\frac{N_{i,v}}{c_i}$  can be retrieved from the decorations of the trees of the transversal sections, except the order N in z of f when the transversal section is not divisible by the variables  $x_i$ .

*Proof.* The proof relies on the fact that the global decorations are given in terms of the equations of the faces of the Newton diagrams. If the vertex we are considering corresponds to a face of the Newton diagram of f, we have the following cases:

- (1) By projection on  $\{x_i = 0\}$  the face is parallel to the z-axis: In this case, the vertex is the first vertex on the top of the tree and its decoration is  $((0, \ldots, N_i, \ldots, 0))$ . The vertex will disappear, but the decorations are 0.
- (2) By projection the face  $\Gamma$  gives a face of the Newton diagram of the transversal section. Either this face is different from the projections of faces intersecting  $\Gamma$ , and it will give a vertex with the decoration  $((N_1, \ldots, N_d))$ , or it is the same face as one of the projections, then the vertex will disappear, but the vertex representing the face have the desired decorations.

Now, if we are on a Newton diagram which appeared in the Newton process, we have again different cases:

- (a) The Newton process is a linear change of variable for z. In this case there are two consecutive vertices which appear with the same  $N_j$ ,  $j \neq i$ . Only one appears in the transversal section, but it bears the decoration we are interested in.
- (b) The Newton process is a Newton map for the transversal section, with the same roots. In this case, either we find the vertex in the transversal sections, or it disappears because there are two faces which project on the same. In any case the remaining vertex gives the decoration we need.
- (c) The Newton process is a Newton map for the transversal sections with more roots. This case is the same as the case above but we have many copies of the Newton diagram. Then the decoration appears many times.

We can notice that not only we recover the decorations of the tree of f, but no new decorations appear in the transversal sections.

Remark 5.14. A particular useful case of the preceding proposition is that all the  $N_i$ 's can be retrieved using the germs of curves obtained by setting all variables  $x_j$  except one as constant. This is a fundamental fact in the proof of the monodromy conjecture for quasi-ordinary polynomial [4].

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